MINOS

A LARGE-SCALE NONLINEAR PROGRAMMING SYSTEM

(For Problems With Linear Constraints)
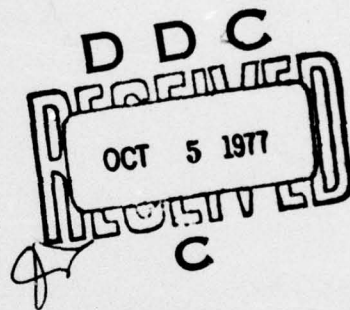
USER'S GUIDE

BY

BRUCE A. MURTAGH AND MICHAEL A. SAUNDERS

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OCT 5 1977

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TECHNICAL REPORT SOL 77-9

FEBRUARY 1977

Systems Optimization Laboratory

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by

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12 132 p.

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SYSTEMS OPTIMIZATION LABORATORY

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15 N00014-75-C-0865,
EY-76-S-03-0326

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I. Introduction

I.1. The Buzz Words

MINOS stands for "a Modular In-core Nonlinear Optimization System." It sounds roughly like the first character in "$-\infty$" (minus infinity), and its purpose is to optimize your own special objective function $f(x)$ by finding a point x which makes $f(x)$ as close to $\pm\infty$ as possible.

Since it is not normally meaningful to go quite that far, MINOS allows you to restrict the variables x to some feasible region specified by a set of linear constraints $(Ax \gtreqless b)$ and a set of upper and lower bounds $(\ell \leq x \leq u)$. In particular, the matrix A may be large and sparse. Of course A may be small and sparse, or even small and dense; in general, the smaller and sparser the better.

Ideally your objective function $f(x)$ should be smooth, the nonlinear part of it should involve as few of the components of x as possible, and you should know how to compute the gradient vector $g(x) = (\partial f/\partial x_j)$ at any feasible point x.

With the important words thus introduced, we will now leave an example problem to speak for itself. The remainder of this section contains the following:

--a statement of the problem;

--a Fortran subroutine to compute $f(x)$ and $g(x)$;

--input cards to be read by MINOS;

--the output produced during solution of the problem.

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The reader should not expect to understand all items immediately. However, the example should help to determine if a particular problem is of the requisite form, and may serve as a useful reference thereafter. Note that "x" in this example is the vector (x, y, z).

Example 1

Objective:	minimize $f(x, y, z) = x^2 + (y - 1)^4 - 4x + z$
Linear Constraints:	$x + y \leq 2$
	$x + y - z = 0$
Bounds:	$x \geq 0, \quad 1 \leq z \leq 5$
Nonlinear Variables:	x and y

Fortran Subroutine to Compute the Nonlinear Part of $f(x, y, z)$

and its Gradient

```
      SUBROUTINE CALCFG( MODE,N,X,F,G,NSTATE,NPROB )
      IMPLICIT   REAL*8(A-H,O-Z)
      DIMENSION   X(N),G(N)
C
C
C     EVALUATE FUNCTION  F   AND GRADIENT  G  AT CURRENT POINT
C     THE INTEGER  N  WILL BE  2  IN THIS EXAMPLE.
C
      T    = X(2) - 1.0
      F    = X(1)**2 + T**4
      G(1) = 2.0*X(1)
      G(2) = 4.0*T**3
      RETURN
      END
```

(The parameters of subroutine CALCFG are defined in Section III.3.)

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Input Cards to be Read by MINOS (SPECS data, followed by MPS data)


```
BEGIN SPECS FOR EXAMPLE PROBLEM
     MINIMIZE
     OBJECTIVE              OBJ.ROW
     RHS                    RHS1
     BOUNDS                 BND1

     ROWS                        10    * These 3 numbers must be
     COLUMNS                     20    * over-estimates of the dimensions
     ELEMENTS                   100    * of the constraint matrix

     ITERATIONS                 100
     SOLUTION                   YES

     PROBLEM NO.                  1
     NONLINEAR VARIABLES          2
     SUPERBASICS                  3
     LINESEARCH TOLERANCE       0.1
END EXAMPLE SPECS
```

```
NAME            EXAMPLE1
ROWS
 N  OBJ.ROW
 L  ROW001
 E  ROW002
COLUMNS
    X           OBJ.ROW        -4.0     ROW001      1.0
    X           ROW002          1.0
    Y           ROW001          1.0     ROW002      1.0
    Z           OBJ.ROW         1.0     ROW002     -1.0
RHS
    RHS1        ROW001          2.0
BOUNDS
 FR BND1        Y
 LO BND1        Z               1.0
 UP BND1        Z               5.0
ENDATA
```

Column numbers for fields on MPS data cards:

		NAME		NAME		VALUE			NAME		VALUE	
12	5	12	15	22	25		36	40	47	50	61

NOTE: The nonlinear variables X and Y must appear _first_

in the COLUMNS section of the MPS data.

Output from MINOS

```
M I N O S  ---  VERSION 3.1   FEB 1977
= = = = =
```

PROBLEM SPECIFICATIONS

```
BEGIN SPECS FOR EXAMPLE PROBLEM
   MINIMIZE
   OBJECTIVE             OBJ.ROW
   RHS                   RHS1
   BOUNDS                BND1

   ROWS           10      * THESE 3 NUMBERS MUST BE
   COLUMNS        20      ** OVER-ESTIMATES OF THE DIMENSIONS
   ELEMENTS      100      * OF THE CONSTRAINT MATRIX

   ITERATIONS    100
   SOLUTION      YES

   PROBLEM NO.           1
   NONLINEAR VARIABLES   2
   SUPERBASICS           3
   LINESEARCH TOLERANCE  0.1
END EXAMPLE SPECS
```

PARAMETERS

```
MPS INPUT DATA.
ROW LIMIT.....:    10          COLUMN LIMIT...:    20          AIJ TOLERANCE..  1.00E-10
LISTING LIMIT..    0           ERROR MSG LIMIT     10          UPP BND DEFAULT  1.00E+30
                                                LOW BND DEFAULT  0.0

FILES.
INPUT MPS DATA.    5           SOLUTION FILE..     0           OLD BASIS MAP..     0          NEW BASIS MAP..    0
INSERT........     0           PUNCH..........     0           LOAD NAMES.....     0          DUMP NAMES.....    0
(SCRATCH).....     8

FREQUENCIES.
CHECK ROW ERROR   25           FACTOR (INVERT)     50          SAVE BASIS MAP.    100          LOG ITERATIONS.    1

LP PARAMETERS.
ITERATIONS....:   100          PARTIAL PRICE..     1           MULTIPLE PRICE.     0          WEIGHT ON OBJ..    0.0
TARGET OBJ.VAL.  0.0           ALIGNMENT TOL..  1.00000

NONLINEAR PROBLEMS.
PROBLEM NO....:    1           NONLINEAR FCNS.     2           NONLINEAR FCNS.     1          DERIVATIVE LVL.    1
SUPERBASICS...:    3           HESSIAN SIZE...     3           CONJ-GRD MODE..     1          VERIFY.........    1
LINESEARCH TOL.  0.10000       REDUCED-GRD TOL  0.20000

MISCELLANEOUS.
PRINT LEVEL...:    0           DEBUG LEVEL....     0           LU FACTOR.......   YES
LU ROW TOL....  0.00100        LU COL TOL.....  0.10000        LU MOD TOL.....  0.99000
```

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INPUT LISTING
----
  1   NAME
  2   ROWS
  6   COLUMNS
 11   RHS
 13   BOUNDS
 17   ENDATA

EXAMPLE1

NAMES SELECTED

OBJECTIVE: OBJ.ROW   (MIN)   1
RHS: RHS1   1
RANGES: 0
BOUNDS: BND1   3

MATRIX STATISTICS

	TOTAL	NORMAL	FREE	FIXED	BOUNDED
ROWS:	3	1	1	1	0
COLUMNS:	4	1	1	1	1

NO. OF MATRIX ELEMENTS: 8   DENSITY: 66.667
NO. OF REJECTED COEFFS: 0   AIJTOL: 1.00000E-10   (EXCLUDING OBJ AND RHS)
BIGGEST AND SMALLEST COEFFS: 1.00000E+00   1.00000E+00

LENGTH OF HASH TABLE, COLLISIONS: 101   0

NO. OF INITIAL BOUNDS PROCESSED: 0
NO. OF SUPERBASICS SPECIFIED: 0

PARTIAL PRICE FACTOR, PARTITION: 1   3

NUMBER OF WORDS OF CORE AVAILABLE: 11776

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ITERATIONS

CRASH OPTION   1
FREE LOGICALS AND STRUCTURALS IN BASIS: 1   1

```
FACTORIZE   1    DEMAND    0    ITERATION   0    INFEAS   0    OBJECTV   0.0
SLACKS      2    LINEAR    0    NONLINEAR   0    ELEMS    4    DENSITY   44.4
PH BUMPS    0    SPIKES    0    CORE REQD  153   L LIMIT  2818  U LIMIT   56.36
LU BUMPS    0    SPIKES    0    AIJ ELEMS   2    L ELEMS   2    U ELEMS   1    F ELEMS   1
ITN    0 -- INFEASIBLE.   NUM =   1    SUM =   1.0000000000D+00                     0   0.0

ITN PH PP NOPT DJ/RG   +SBS -SBS -BS   STEP     PIVOT    NSPY  L  U  NINF SINF/OBJECTIVE   NFG NSB P1M H-COND COND
 1  1  1    2 -1.00+00    3    3   7  1.00+00  -1.00+00    0   2  1   1  1.00000000D+00    0   0   0  0.0     FFFT

ITN   1 -- FEASIBLE SOLUTION.   OBJECTIVE =   3.000000000D+00

 2  3  1    2 -8.00+00    0    0   0  2.0D-01   0.0        1   2  2   0 -1.70975950D+00    4   1   + 1.00+00  FFTF
 3  4  1    2  2.00-06    0    0   0  8.0D-01   0.0        1   2  2   0 -1.71026580D+00    6   1   + 1.00+00  TTTT
 4  3  1    1 -1.80-01    6    0   0  1.00+00   0.0        1   2  2   0 -1.72423930D+00    8   2   + 4.60+00  TTTF

RG TOLS REDUCED.   TOLRG =   7.071D-07

 5  1  1    0  1.50-03    0    0   0  1.00+00   0.0        1   2  2   0 -1.72470200D+00    9   2   + 5.00+00  FFFF
 6  4  1    0  1.60-05    0    0   0  1.00+00   0.0        1   2  2   0 -1.72470390D+00   10   2   + 4.90+00  FFFF
 7  4  1    0  2.20-07    0    0   0  1.00+00   0.0        1   2  2   0 -1.72470390D+00   11   2   + 4.80+00  FFTF
 8  4  1    0  1.50-09    0    0   0  1.00+00   0.0        1   2  2   0 -1.72470390D+00   12   2   + 4.80+00  TTTT

BIGGEST DJ =   0.0          NORM RG =   1.461D-09   NORM PI =   7.071D-01

EXIT -- OPTIMAL SOLUTION FOUND.

NO. OF ITERATIONS:              8        PHASE:  3        OBJECTIVE:   -1.72247039371D+00

NO. OF FUNCTION AND GRADIENT CALLS:   12
```

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PROBLEM NAME: EXAMPLE1   OBJECTIVE VALUE: -1.72247039370+00

STATUS: OPTIMAL SOLN   PHASE: 3   ITERATION: 8

OBJECTIVE: OBJ.ROW (MIN)
RHS: RHS1
RANGES:
BOUNDS: BND1

SECTION 1 - ROWS

NUMBER	...ROW..	AT	...ACTIVITY...	SLACK ACTIVITY	..LOWER LIMIT.	..UPPER LIMIT.	.DUAL ACTIVITY	..J
5	OBJ.ROW	BS	-4.12996	4.12996	NONE	NONE	-1.00000	1
6	ROW001	SBS	1.87004	0.12996	NONE	2.00000	0.00000	2
7	ROW002	EQ	0.0	0.0	0.0	0.0	-1.00000	3

SECTION 2 - COLUMNS

NUMBER	.COLUMN.	AT	...ACTIVITY...	.OBJ GRADIENT.	..LOWER LIMIT.	..UPPER LIMIT.	.REDUCED COST.	..J
1	X	SBS	1.50000	-1.00000	0.0	NONE	0.0	4
2	Y	BS	0.37004	-1.00000	NONE	NONE	0.0	5
3	Z	BS	1.87004	1.00000	1.00000	5.00000	0.0	6
4	RHS1	EQ	-1.00000	0.0	-1.00000	-1.00000	0.0	7

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I.2. <u>Formal Introduction</u>

MINOS is a computer program designed to minimize a linear or nonlinear function subject to linear constraints.[†] In algebraic terms it is designed to solve problems of the form

P:
$$\text{minimize} \qquad f(x) + c^T x$$

$$\text{subject to} \qquad Ax \gtreqless b, \qquad \ell \leq x \leq u$$

given that $f(x)$ is a continuous differentiable function with gradient $\nabla f(x) = [\partial f / \partial x_j] = g(x)$.

Such problems range from small, unconstrained optimization problems through constrained linear least squares (data-fitting) models to large-scale linear programs with (or without) nonlinear terms in the objective. In general the constraint matrix A is assumed to be large and sparse.

Fundamental to the system is an efficient and reliable implementation of the revised simplex method for linear programming (LP).[‡] This combines established sparse-matrix technology with stable numerical methods for computing and modifying a triangular factorization of the usual square basis matrix B.[‡] The code is intended to be suitable for production runs on purely linear problems. The data formats used for various files therefore follow those adopted by current commercial mathematical programming systems. In particular, the

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[†]The theory of the solution technique is described in a companion paper, Murtagh and Saunders (1976).

[‡]Dantzig (1963).

[‡]Bartels and Golub (1969), Bartels (1971), Saunders (1976).

quantities A, b, c, ℓ and u are defined in the format accepted by the IBM Mathematical Programming Systems MPS/360, MPSX and MPSX/370. Ideally the reader of this manual should have experience in using such systems and should (a fortiori) be familiar with the concepts and terminology associated with the simplex method.

Recall that the <u>basic</u> variables associated with B may take solution values anywhere between their upper and lower bounds, and that the remaining variables are defined to be <u>nonbasic</u> at one or other of their bounds. The simplex method always works with solutions of this form. In order to extend the simplex method to problem P we introduce a new class of variables, which will be called <u>superbasic</u>. Both basic and superbasic variables may vary between their bounds, but in the reduced-gradient[†] approach used here their roles are different. Suppose there are ns superbasic variables at some stage. These are essentially free to move in any desirable direction, namely one which will improve the value of the objective functions; in effect they provide the driving force. The basic variables can then be adjusted so that the full set of variables x remains feasible with respect to the linear constraints. If it appears that no improvement can be made with the current set of superbasics, one (or more) of the nonbasic variables is selected to become superbasic, ns is increased and the process is repeated. At all stages, if a basic or superbasic variable encounters one of its bounds, an adjustment occurs in which that variable is made nonbasic and the total number of superbasics is reduced by 1.

[†] Wolfe (1962, 1967).

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Users familiar with LP may interpret the simplex method as being exactly the above process, with ns oscillating between 0 and 1.

At any given stage, the number of superbasics (ns) measures in some sense how far the current point x is from a vertex of the simplex (polyhedron) defining the feasible region. In other words, x lies somewhere on a face of the polyhedron, which defines a sub-space of dimension ns. An important part of MINOS is a stable implementation of a quasi-Newton (or variable-metric) method for optimizing the superbasic variables within each appropriate subspace. This method uses a triangular matrix of dimension ns to approximate the reduced Hessian (i.e., a suitably transformed sub-section of the matrix of second derivatives, $[\partial^2 f/\partial x_i \partial x_j]$). Good rates of convergence are usually achieved.

If the number of superbasic variables ever grows as large as 100 or 200 the storage required for the quasi-Newton procedure starts to become excessive. In such cases MINOS will commence using a conjugate-gradient method,[‡] which requires very little storage. The rate of convergence normally drops severely, but at present this is the only practical alternative for problems with very many nonlinearities.

The concept of superbasic variables can be valuable even for purely linear problems. As an example, the solution to a particular problem may provide a feasible but non-vertex starting point for a modified version of the problem (e.g., if the bounds were altered on some nonbasic variables). The re-starting features in MINOS allow

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[†]Davidon (1959), Fletcher and Powell (1963).
[‡]Fletcher and Reeves (1964).

the user to take advantage of such situations easily. On a less routine level, MINOS also provides a framework within which experiments can be performed on large problems using non-simplex steps, perhaps with a view to reducing the total number of iterations required.

I.3. The Canonical Form

Without loss of generality, problem P may be re-cast in the form

$$\text{minimize} \quad f(x_N) - s_{obj}$$

$\bar{P}$:

$$\text{subject to} \quad [A_N \ A_L \ b \ I] \begin{bmatrix} x_N \\ x_L \\ \rho \\ s \end{bmatrix} = 0, \qquad \ell \leq \begin{bmatrix} x_N \\ x_L \\ \rho \\ s \end{bmatrix} \leq u$$

$$\text{given} \quad \nabla f(x_N) = g(x_N)$$

where the following definitions hold:

x_N = the <u>nonlinear variables</u>, i.e., those that are directly involved in the function $f(x_N)$.

x_L = the <u>linear variables</u>, i.e., the remaining part of x.

ρ = the <u>right-hand side</u> variable, which has bounds $\ell_\rho = u_\rho = -1.0$

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s = the <u>slack</u> or <u>logical</u> <u>variables</u> (one for each row of A).

s_{obj} = the slack variable corresponding to the linear part (if any) of the objective function. The term $c^T x$ in problem P is thus imbedded in the constraint matrix as the "obj"-th constraint, $c^T x + s_{obj} = 0$, with s_{obj} being a free variable. (This row may be omitted if $c = 0$ or if c is included in f.)

A = $[A_N \ A_L]$

m = the number of rows in A.

n = the number of columns in A.

nn = the number of nonlinear variables in x_n.

ns = the number of superbasic variables.

At any particular stage, the columns of $[A \ b \ I]$ are implicitly ordered as follows:

$$[A \ b \ I]\hat{P} = \begin{array}{ccc} m & ns & n+1-ns \\ \boxed{B} & \boxed{S} & \boxed{N} \end{array}$$

Basics, as in
simplex method Superbasics

Nonbasics
(variables at upper
or lower bound)

where $\hat{P}$ is a permutation. The nonlinear variables may end up anywhere in B, S or N, and similarly the superbasic variables in S may be linear or nonlinear. The only requirement is that the basis matrix B be nonsingular.

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The above partitioning should be remembered when interpreting
output from MINOS and when modifying certain basis files for input.

II.   Nonlinear Problems

At any given stage, the reduced-gradient method regards the objective as an essentially unconstrained function of the super-basic variables alone. (Basic and nonbasic variables are eliminated by means of the active constraints.) It is therefore possible to use unconstrained optimization methods on some "reduced function" of the ns superbasic variables. Suppose that the original objective has gradient g and Hessian G (i.e., the matrix of second partial derivatives). The gradient and Hessian for the reduced function are then $Z^T g$ and $Z^T G Z$, where Z is defined in terms of the current basis as the $n \times ns$ matrix

$$
Z = \begin{bmatrix} -B^{-1}S \\ I \\ O \end{bmatrix} \begin{matrix} m \\ ns \\ n+1-m-ns \end{matrix} \qquad .
$$

$$ns$$

Note that it would not be practical to compute Z or $Z^T G Z$ (assuming m to be large and ns to be more than 1 or 2). Any unconstrained optimization method to be used must therefore make do with vectors of the form $Z^T g$ or Zp, which can be computed with relative ease.

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II.1 The Quasi-Newton Method

Whenever storage permits, as specified by a HESSIAN DIMENSION
card in the SPECS file (see Section IV), MINOS maintains a quasi-
Newton approximation to the reduced Hessian thus:

$$R^TR \approx Z^TGZ$$

where R is a dense, nonsingular, upper triangular matrix of
dimension ns. This normally provides superlinear convergence to
a minimum within the subspace defined by each set of superbasic
variables. The factorized form R^TR is justified by the fact
that Z^TGZ must be positive definite (or at least semi-definite)
at such a minimum. It also guarantees that the equations

$$R^TRs = -Z^Tg, \qquad p = Zs$$

will generate a direction p along which the original objective
will decrease (since the projection of p along the gradient is
$p^Tg = s^TZ^Tg = -s^TR^TRs = -\|Rs\|_2^2 < 0$).

Various modifications are made to R each iteration to
account for new curvature information and to allow for changes in
the number of superbasics. For example, when a nonzero step is
taken, a new matrix $\bar{R}$ is computed in two stages to satisfy

$$\bar{R}^T\bar{R} = R^TR + vv^T - ww^T$$

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for some vectors v and w (which are determined by the so-called "Complementary DFP" or "BFGS" formula, e.g., see Broyden (1972)). The result of such operations on R is briefly indicated in the iteration log by two integers labelled "R1M" (short for Rank-one Modifications). In general, these numbers are:

positive if the modifications were successful;

zero   if the change in reduced gradient was too small for v

          and w to be meaningful (in which case R is unaltered);

negative if numerical error prevented a particular modification.

Another useful quantity appearing in the iteration log is an estimate of the "condition number" of $R^T R$. See the discussion of H-CONDN in Section VII ; also see Section II.3.

II.2. The Conjugate-Gradient Method

If insufficient storage has been allocated for R, MINOS will discontinue the quasi-Newton method for generating search directions and switch to a conjugate-gradient method (which does not need R). The dimension of R may be specified separately from the upper limit on superbasics, thus:

SUPERBASICS   200

HESSIAN DIMENSION   100

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In this example a conjugate-gradient method will be used if the number of superbasics exceeds 100. Quasi-Newton iterations can be suppressed completely by the card

<div align="center">HESSIAN DIMENSION   0</div>

(The total no. of iterations will probably increase significantly.)

Five versions of "CG" are currently available via the SPECS card

<div align="center">CONJUGATE GRADIENT VERSION   k</div>

for the following values of k:

k	Method
1	Fletcher and Reeves (1964)
2	Polak and Ribiere (1969)
3	Perry (1976)
4	Memoryless DFP
5	Memoryless Complementary DFP

The general form of the first three methods is

$$s = -Z^T g + \beta s_{old}$$

<div align="center">17</div>

where s and s_{old} are the new and previous search directions for the superbasics, and β depends on the method. The other two methods are derived from the corresponding quasi-Newton formulas, taking the initial reduced Hessian approximation to be I. They add to s a term of the form

$$\gamma(Z^Tg - Z^Tg_{old}) \ .$$

Restarts $(\beta = \gamma = 0)$ occur at every change to the basic or super-basic sets, or after ns consecutive unconstrained steps with a particular set of ns superbasics.

The state of research on conjugate-gradient methods is such that no clear recommendation can be made for the choice of k. If an accurate linesearch is used, all methods are rather similar and $k = 1$ or 2 should be tried. If the objective function is so expensive to evaluate that an inaccurate search is necessary, $k = 3$ or 5 may be preferable. Method 4 appears from limited experience to be the least successful. Considerable further research is needed in this area.

II.3. Scaling Nonlinear Variables

No rigid rules can be given here that would apply usefully to all nonlinear functions, but one quantity in particular may be worth examining if it seems that optimization is proceeding very slowly. This quantity is $\text{cond}(G)$, **the condition number of the Hessian**, which may be very sensitive to scaling.

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We shall illustrate with a particular function of two variables, viz.

$$f(x, y) = x^b y^c$$

where b and c are known parameters. It is readily verified that the Hessian for f is

$$G(x, y) = \begin{bmatrix} \partial^2 f/\partial x^2 & \partial^2 f/\partial x \partial y \\ \partial^2 f/\partial x \partial y & \partial^2 f/\partial y^2 \end{bmatrix} = \frac{f}{x^2 y^2} \begin{bmatrix} b(b-1)y^2 & bcxy \\ bcxy & c(c-1)x^2 \end{bmatrix} .$$

Furthermore, G may be factorized as

$$G(x, y) = \frac{f}{x^2 y^2} LL^T$$

where

$$L = \begin{bmatrix} \beta y & \\ \gamma x & \delta x \end{bmatrix}$$

$$\beta^2 = b(b-1)$$

$$\gamma = bc/\beta$$

$$\delta^2 = c(1 - b - c)/(b - 1).$$

A triangular factorization of this kind provides a lower bound on the condition number of G, in terms of the ratio of the largest and smallest diagonals of L. In this case, suppose that $|\beta y| \geq |\delta x|$ at the optimal solution. Then

$$\text{cond}(G) = \text{cond}(L)^2 \geq \left(\frac{\beta y}{\delta x}\right)^2 .$$

Assuming that β and δ are reasonable numbers (of order 1), it appears that the ratio y^2/x^2 at the optimal solution may be a useful indicator of the sensitivity of the optimum to small changes in the data. To make the estimate of $\text{cond}(G)$ close to 1, x and y should be scaled to have approximately the same values at the solution. (The same conclusion would be drawn if $|\beta y| \leq |\delta x|$.)

Warning: In the reduced-gradient method, scaling for the above reasons should not be at the expense of upsetting the scale of the constraint matrix, A. This is because the quantity of interest is really $\text{cond}(Z^T G Z)$ rather than $\text{cond}(G)$, where Z is a matrix whose scale (and condition) depend on A.

Note that if certain variables (and hence columns of A) are scaled to improve $\text{cond}(G)$, it may be necessary to scale certain rows of A to maintain the condition of A and $Z^T G Z$. It may then be necessary to scale certain other columns of A, and so on. Clearly, it is preferable to be careful with scaling during the initial formulation of a problem.

Finally, the analysis suggested above will usually be applicable to certain obvious rows and columns of G, rather than to G as a whole; for example, if G is block-diagonal, each block may be analyzed in turn.

III. Data Required from the User

The following three items must be supplied by the user:

1. the SPECS file   to specify certain run-time parameters;
2. the MPS file   to specify the linear constraints in
                           standard MPS format;
3. subroutine CALCFG   to compute the nonlinear part of the
                           objective function and its gradient.

The user may also (optionally) supply

4. certain BASIS files   to define an initial solution.

Since several or all of the various input files could be part of the
card input stream, it is important to note that MINOS reads data in
the following order:

(a) SPECS data

(b) MPS data

(c) BASIS data

(d) any data required by subroutine CALCFG.

III.1. The SPECS File

Certain parameters may be set by the user via a list of
"problem specifications," hereafter called SPECS. These are assumed
to be a deck of 80-character card images of the form

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BEGIN SPECS FOR FRED

⋮

(list of keywords and values)

⋮

END FRED

where the words BEGIN and END are themselves keywords. (Other charac-
ters on BEGIN and END cards are ignored.) The SPECS file must exist;
it is normally the first data set in the card input stream (file 5).

Broadly speaking, each card in the SPECS file contains a
sequence of items punched in free format (i.e., separated by at least
one blank or =). The items selected from each card are

1. The first word   (the keyword; only the first 3 characters are
significant);

2. The second word   (8 characters, treated as two sets of 4);

3. The first number (up to 8 characters; either integer or real).

In the following examples the significant characters are underlined:

            OBJECTIVE   PROFIT

            ROWS   500

            AIJTOL   1.0E-6

            LOWER BOUND   -1.0

            SOLUTION FILE   20

            OLD BIT MAP ON FILE 9

Only occasionally is the second word meaningful (e.g., PROFIT and FILE above). Words such as BOUND and BIT MAP are extracted but later ignored.

Blank cards are allowed, and comments (i.e., any characters) may occur after an asterisk, e.g.,

NONLINEAR VARIABLES   20   * X005, X010, ..., X100.

Default Values

Parameters that are not specified in the SPECS file assume certain default values, most of which should be acceptable for normal use. In the following list the default values are given for all legal keywords. Where possible (except for some of the items printed in lower case) the list itself constitutes a valid SPECS file. Note, however, that its effect would be essentially the same as the file

BEGIN

END

(which is also valid).

```
BEGIN   SPECS FILE SHOWING DEFAULT VALUES FOR KEYWORD PARAMETERS
*
*
*   KEYWORDS FOR MPS DATA
*
*
    MINIMIZE
    OBJECTIVE   =    the first name encountered
    RHS         =              ditto
    RANGES      =              ditto
    BOUNDS      =              ditto

    ROWS                         100     *
    COLUMNS                      300     *  (or  3xROWS)
    ELEMENTS (COEFFICIENTS)     1500     *  (or  5xCOLS)
    AIJ ZERO TOLERANCE        1.0E-10    *
    LOWER BOUND                  0.0     *
    UPPER BOUND               1.0E+30    *
                                         *
    INPUT FILE                     5     *  (MPS data)
    LIST LIMIT                     0     *  (for printing of MPS data)
    ERROR MESSAGE LIMIT           10     *  (during MPS input)
*
*
*
*   KEYWORDS FOR THE SIMPLEX ALGORITHM
*
*
    CRASH OPTION                   1     *  (Triangular CRASH on all vars)
    ITERATIONS                   300     *  (3xROWS + 10xNONLINEARS)
    PARTIAL PRICE                  1     *
    MULTIPLE PRICE                 0     *  (BEWARE - not like normal LP)
                                         *
    WEIGHT ON OBJECTIVE          0.0     *  (during Phase 1)
                                         *
    FACTORIZATION FREQUENCY       50     *  (i.e. INVERT FREQUENCY)
    CHECK FREQ                    25     *  (test on row residuals)
    SAVE  FREQ                   100     *  (for BIT MAP)
    LOG   FREQ                     1     *  (for ITERATION LOG)

    SOLUTION    IF OPTIMAL, INFEASIBLE, OR UNBOUNDED
*
*
*
*   FILES FOR SAVING AND LOADING BASIS DESCRIPTIONS
*
*
    OLD BASIS FILE                 0     *  (BIT MAP)
    NEW BASIS FILE                 0     *  (BIT MAP)
    INSERT    FILE                 0     *  (Industry standard)
    PUNCH     FILE                 0     *  (for saving of INSERT data)
    LOAD      FILE                 0     *  (NAMES and VALUES)
    DUMP      FILE                 0     *  (for saving of  LOAD  data)
    SOLUTION  FILE                 0     *  (Don't confuse with the
                                         *   printed solution)
```

```
*
*
*
*    PARAMETERS FOR NONLINEAR PROBLEMS
*
*

     PROBLEM NUMBER                  0       *   (Parameter NPROB for CALCFG)
     NONLINEAR VARIABLES             0       *   (Default is pure LP)
     SUPERBASICS LIMIT               1       *   (or  NONLINEARS + 1)
     HESSIAN DIMENSION               1       *   (min( SUPERBASICS, 30 ))
     CONJUGATE GRADIENT VERSION      1       *   (FLETCHER-REEVES)
                                             *
     VERIFY                         NO       *   (Suppress gradient check)
                                             *
     LINESEARCH TOLERANCE          0.01      *   (Must be between 0.0 and 1.0)
     REDUCED-GRADIENT TOL          0.2       *   (ditto)
*
*
*
*    MISCELLANEOUS
*
*

     DEBUG LEVEL                     0       *
     PRINT LEVEL                     0       *
     IMBED                         YES       *   (Store part of LU file in  A)
                                             *
     LU ROW TOLERANCE             0.001      *   (for initial LU factorization)
     LU COL TOL                    0.1       *   (ditto)
     LU MOD TOL                    0.99      *   (for updating LU factors)
*
*
*
*    THE FOLLOWING KEYWORDS EACH ENABLE A CERTAIN OPTION.
*    THEY SHOULD NOT APPEAR IF THE OPTION IS NOT WANTED.
*
*
     ALIGNMENT TOLERANCE           1.0       *   (Experimental, for pure LP)
     CALL CALCFG IF OPTIMAL                  *   (One last entry with NSTATE=2)
     TARGET OBJECTIVE VALUE        0.0       *   (Experimental, for pure LP)

END LIST OF DEFAULT VALUES


NOTE:  The keywords  BEGIN  and  END  are required.
```

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III.2. The MPS File

After the SPECS file has been successfully read, the data
specifying the constraint matrix is input in standard MPS format.
This format is defined under the title "CONVERT DATA" in IBM
document number SH20-0968-1, "Mathematical Programming System-Extended
(MPSX), and Generalized Upper Bounding (GUB)," pp. 199-209.
Alternatively, see IBM document number H20-0476-2, "Mathematical
Programming System/360 Version 2, Linear and Separable Programming --
User's Manual," pp. 141-151.

The various sections of the data must be in the following
order:

> NAME
> ROWS
> COLUMNS
> RHS
> RANGES
> BOUNDS
> ENDATA

Notes on MPS Data

1. The relevant column numbers and contents for the fields on each
   data card are as follows:

Field 1	Field 2	Field 3	Field 4	Field 5	Field 6
2-3	5-12	15-22	25-36	40-47	50-61
code	name	name	value	name	value

2. The NAME card may define a name of any length, beginning in column 15. Only the first 8 characters are saved for printing with the solution.

3. The ROWS section may define constraints to be of type $=$, $\geq$, $\leq$ or free, using the characters E, G, L or N respectively (in column 2 or column 3). It is not essential to have an objective row; thus there may be 0, 1 or more rows of type N.

4. The COLUMNS section may define 1 or 2 matrix elements per card. For nonlinear columns, it is also acceptable to have 0 elements (i.e., just the name of the variable).

   NOTE: Nonlinear variables must occur <u>first</u> in the COLUMNS section, ordered in a manner that is consistent with the array X in the user-written subroutine CALCFG.

5. The RHS section may contain several RHS's. A specific one may be requested via the specifications, otherwise the first RHS is taken.

6. Similarly for the RANGES section.

7. Similarly for BOUNDS. The types of bound allowed are UP, LO, FX, FR, MI, PL.

   <u>Warning</u>: If a variable is essentially unconstrained the bound type FR (free) should always be used. <u>Never</u> use a large negative LO bound. Similarly if a variable XI has bounds $-\infty \leq XI \leq 5$ (say), then 2 cards are required as follows:

   ```
   FR BND1      XI
   UP BND1      XI            5.0
   ```

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8. The INITIAL bounds set. The name INITIAL is reserved to specify (optionally) a special bounds set which can be used to assign initial values to nonlinear variables. The INITIAL bounds set must appear after any normal bounds sets (if any). A warning is given if it is the first set encountered after the BOUNDS card. The following example illustrates the bound types allowed:

Field 1	Field 2	Field 3	Field 4
FX	INITIAL	XN1	10.0
LO	INITIAL	XN2	
UP	INITIAL	XN3	

In this example:

(a) The nonlinear variable XN1 will be initialized at value 10.0. (Its initial state will be superbasic.) If 10.0 lies outside the bounds on XN1, the initial value will be the nearest bound.

(b) Nonlinear variables XN2 and XN3 will be initialized at their lower and upper bound values respectively. (Their initial state will be nonbasic.)

(c) Any nonlinear variables that are not specified directly will initially be nonbasic at their smallest bound (in absolute value).

Note: The INITIAL bounds set provides an important facility if used in conjunction with CRASH OPTION 2 (which prevents nonlinear variables from being selected for the initial basis). By construction, all

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nonlinear variables will then be either superbasic or nonbasic, and they will be held fixed at their initial values while the remaining (linear) variables are optimized by a normal Phase 1/Phase 2 simplex procedure. Once it has been determined that the associated linear program is either optimal or infeasible, the nonlinear variables will be allowed to vary and optimization will continue in normal reduced-gradient mode.

It follows that an INITIAL bounds set and CRASH OPTION 2 should be used if good starting values are known for all of the non-linear variables. (At least one FX indicator must be used to achieve the above effect, and the SUPERBASICS card should specify a suitably large number. If FX is used for all nonlinear variables, then CRASH OPTION 2 may be omitted.)

If a starting point is specified by any other means (i.e., via OLD BASIS, INSERT BASIS or LOAD BASIS cards), the INITIAL bounds set will be ignored. This allows the user to re-start a run without altering the MPS file.

Warnings, provisos, etc.

1. The constraint matrix is not scaled by the program, so the user should try to ensure that

    a) all elements a_{ij} are reasonably close to 1.0;

    b) units are chosen so that all x_j satisfy $10^{-2} \leq |x_j| \leq 10^2$ or $x_j = 0$.

    These are by no means absolute requirements but are recommended for numerical reliability (and to ensure that there are at least some significant figures in the printed solution).

2. The program _does_ make allowance for the scale of the objective function and, to a lesser extent, the RHS. The sizes of the objective and b_i elements are therefore somewhat less critical.

3. The nonlinear objective function is evaluated only at feasible points. The BOUNDS section should in general include bounds on all nonlinear variables so that "feasible points" are indeed away from singularities in the objective and its gradient, e.g., if $f(\underset{\sim}{x}) = 1/x_1 x_2^2$ and it is known that $x_i \approx 10$ at the solution, it would be advisable to include bounds forcing $x_i \geq 1$ (say).

III.3. Subroutine CALCFG

Computation of the Objective Function and its Gradient

Let $f(x)$ be the nonlinear part of the objective function, with corresponding gradient vector $g(x)$. Both $f(x)$ and $g(x)$ must be computed by a subroutine called CALCFG, whose specification and parameters are as follows. (Recall that any linear terms in the objective function may be included as part of $f(x)$ or as a free row in the constraint matrix A.)

Specification

```
          SUBROUTINE CALCFG (MODE, N, X, F, G, NSTATE, NPROB)
          IMPLICIT   INTEGER(I-N), REAL(A-H, O-Z)
          DIMENSION  X(N), G(N).

    C
    C     On machines with short word length, e.g., IBM, Univac, it is
    C     necessary to use double precision floating-point.  Hence
    C     IMPLICIT   INTEGER(I-N), REAL*8(A-H, O-Z)
```

MODE   (Input, output) A control parameter defining the type of function and derivative information required.

On entry:

MODE = 2 means the user must assign the function value $f(x)$ to F and the gradient values $g(x)$ to $G(j)$, $j = 1, 2, \ldots, N$, using the values of X stored in $X(j)$, $j = 1, 2, \ldots, N$. (Other values of MODE not yet implemented.)

On exit:

If the user wishes to terminate solution of the current problem, MODE should be given a negative value (viz. -1).

N   (Input) The number of variables involved in $f(x)$.

Note: These must be the first N variables in the constraint matrix A. In general, A will be of dimension m by n with $N < n$.

X   (Input) An array of dimension N containing the current values of the nonlinear variables x.

Note: The contents of X should not be altered, except perhaps if NSTATE = 2 (see below).

F   (Output) The computed value of $f(x)$.

G   (Output) The computed gradient vector $g(x)$. The user must store the partial derivative $\partial f/\partial x_j$ in $G(j)$, $j = 1, 2, \ldots, N$.

NSTATE   (Input)   A control parameter informing the user of the current state of the optimization process.

NSTATE = 0   if there is nothing special about the current entry.

NSTATE = 1   if this is the <u>first</u> entry to CALCFG. The user may wish to read data or set up certain local arrays for use in computing F and G during the present and subsequent entries to CALCFG.

NSTATE = 2   if the current solution in X is optimal. The user may wish to perform additional computations on X.

<u>Notes about entry with NSTATE = 2</u>

1. This will occur only if the user includes a card of the form

        CALL OBJECTIVE WHEN OPTIMAL

    in the problem specifications.

2. This will be the <u>last</u> entry to CALCFG.

3. The entry will not occur if iterations are terminated before optimality is recognized.

4. The entry will be made before the solution is output to the printer and/or the Solution File. Any changes to the nonlinear variables X will therefore appear in the solution output. (For example, one might wish to round the elements of X to the nearest integer.)

NPROB   (Input) · An integer which can be set by the user via a

card of the form

<div style="text-align:center">PROBLEM NUMBER 5</div>

in the problem specifications. This parameter may be

used to branch to one of several function calculations

within CALCFG. The user may ignore it or use and change

it as desired.

III.4. BASIS Files

Four distinct methods are available for loading and saving

basis descriptions.[†] They are invoked by SPECS cards of the following

form. (Just one option per card. The file numbers may be whatever

the user chooses.)

Loading	Saving
OLD BASIS FILE 10	NEW BASIS FILE 11
INSERT FILE 20	PUNCH FILE 21
LOAD FILE 30	DUMP FILE 31
CRASH OPTION 0, 1 or 2	SOLUTION FILE 50

The keywords OLD, INSERT, LOAD are mutually exclusive. If more

than one positive file number is specified, the order of precedence

is that just given. If no starting files are specified, one of the

CRASH options takes precedence.

The keywords NEW, PUNCH, DUMP, SOLUTION may all specify

positive file numbers. Hence from zero to four files may be created

at the end of a run, in the order just given. In addition, a NEW

[†]See also the INITIAL bounds set in Section III.2.

BASIS FILE may be created every k iterations during a run, where k is specified by a SAVE FREQUENCY card.

Normally the various basis files are used in the following circumstances. Fuller details are given in Section V.

1. OLD/NEW bit map.[†] This is the most compact and efficient method for starting from an earlier run on either the same problem or one of the same dimensions.

2. INSERT/PUNCH. This provides compatibility with certain commercial systems (e.g., MPS/360, MPSX/370, MPS III, APEX). The PUNCH file from a particular problem can be used as an INSERT file for restarting solution of the same problem. Occasionally it may be possible to modify the INSERT file and/or the problem and still obtain a useful advanced basis.

   The standard MPS format has been slightly generalized within MINOS to allow saving and reloading of nonbasic solutions.

3. LOAD/DUMP. This is similar in vein to INSERT/PUNCH but allows more direct specification of a list of basic and superbasic variables (both structurals and logicals) and their desired starting values. It is usually easier to ensure that a DUMP file (rather than a PUNCH file) is suitable for restarting a modified problem.

[†]So-called because each variable is represented by just one character.

4. SOLUTION file. This may be used to communicate solution values (e.g., x, π, objective gradient, bounds) to another program for further computation. It uses essentially the same format as that used for the printed solution. For greater precision it uses format 1PE16.6 for all values, rather than F16.5.

IV. Problem Specifications

IV.1. The SPECS file -- Syntax

Each card in the SPECS file must be of the form

⟨keyword⟩ ⟨identifier⟩ ⟨idlist⟩ ⟨number⟩ ⟨comment⟩

where items may be written in free format with one or more blanks or
'='s as separators. In general (depending on ⟨keyword⟩) not all
items need be present.

⟨keyword⟩   A string of non-blank characters beginning anywhere
                 in columns 1 through 70. A keyword is defined by its
                 first 3 characters; any others are ignored.

⟨identifier⟩   A list of 8 consecutive characters, the first of
                 which is any non-blank character except the following,
                 which are hereby defined to be numeric characters:

                 0 1 2 3 4 5 6 7 8 9 + -

                 Valid use of identifiers:

                 OBJECTIVE = COST ROW

                 RHS   = ...Z.

                 RANGES   = RNG001

                 BOUNDS   = LOW1.E-5

                 Invalid use of identifiers:

                 OBJECTIVE = -Z

                 BOUNDS   = 2Q

                 BOUNDS   = 1.0E-7

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Note that blanks may be imbedded in the last 7

characters, but the same blanks must then appear

in the corresponding MPS data cards.

⟨idlist⟩   A list of items separated by blanks or '='s. The

first character of each item is not one of the numeric

characters listed above. All items in the list are

ignored.

Valid use of idlist:

WEIGHT ON OBJ (DURING PHASE ONE) = 100.0

Invalid use of an idlist:

WEIGHT ON OBJ (DURING PHASE 1) = 100.0

⟨number⟩   A list of up to 8 consecutive characters, the first

of which is numeric. The number may be either an

integer or a real constant. It will be treated as an

integer if all remaining characters are numeric; it

should then be typed in I format.

Valid use of integers:

ROWS = 2000

COLS = +2000

Invalid use of integers:

ROWS = 1.0E+3

A number will be treated as a real if any character

other than the first is '.', 'E' or 'D'; it should

then be typed in F, E or D format.

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<u>Valid</u> use of <u>reals</u>:

AIJ TOLERANCE = 0.00001

WEIGHT ON OBJ = +1.0E+5

LOWER BOUND = -1.0D+30

<u>Invalid</u> use of <u>reals</u>:

LOWER BOUND = .1E-5

(comment)   Scanning of a specification card is terminated under

the following conditions:

1. if the first non-blank character is an '*' (i.e.,

    if the first character in the keyword is an '*'),

2. if any blank or '=' is followed by an '*'.

3. If an <u>integer</u> or a <u>real</u> has been recognized.

Hence a comment may be any characters following an

'*' or an <u>integer</u> or a <u>real</u>.

<u>Valid</u> use of comments

*

OBJ   = COST 02 * 2ND ALTERNATIVE

ROWS 1000 (OR LESS)

WEIGHT ON OBJ = 10.0 IN PHASE 1

IV.2.   The SPECS file -- Keywords

The following is an alphabetical list of recognized keywords.

A typical use of each keyword is given, along with a definition of

the quantities involved and various comments or warnings. In many

cases, the value associated with a keyword is denoted by a letter

such as k, and allowable values for k are subsequently defined.

● AIJ TOLERANCE   0.001

During input of the MPS file, matrix coefficients a_{ij} will be
ignored if $|a_{ij}|$ is less than or equal to the specified tolerance.

● ALIGNMENT TOLERANCE   1.0

A parameter invoking a certain algorithmic option which affects
the search direction used for superbasic variables. Specifically,
a subset of the superbasics is selected such that each member x_j
is close to its one of its bounds and has a significant gradient
component moving x_j towards that bound. The search direction is
scaled in a way that will cause all members of the subset to reach
their bounds simultaneously. All members will therefore be "squeezed"
out of the superbasic set, unless a basic variable happens to reach
a bound first.

1. The ALIGNMENT keyword should not appear if alignment is not
    required.

2. Alignment is unlikely to be successful on nonlinear problems.

3. Use of the specified tolerance depends on the particular version
    of subroutine ALIGN.

● BOUNDS   BOUND01

The 8-character name of the bound set to be selected from the MPS file.

1. BNDS is a valid alternative keyword.

2. If no BOUNDS keyword exists, or if the name specified is blank,
    the <u>first</u> bound set in the MPS file will be selected.

3. If the MPS file contains one or more bound sets, but the user
    wishes <u>no</u> bound set to be used, some dummy name should be specified,
    e.g., BOUNDS = NONE.

● CALL OBJECTIVE WHEN OPTIMAL

This requests a final call to the user's subroutine CALCFG when
(and if) an optimal solution has been obtained.

1. Refer to the parameter "NSTATE" in Section III.3.

2. The CALL keyword should not appear if a final call is not
    required.

● CHECK FREQUENCY   k

The residual vector $r = b - s - Ax$ will be computed every k-th
iteration. Refactorization of the current basis B is requested
if $\|D^{-1}r\|$ is too large, where $D = \text{diag}(d_i)$ and d_i is the largest
element in the i-th row of B. Normally k should be about 25 or
half the FACTORIZATION FREQUENCY.

● COEFFICIENTS   3000

See ELEMENTS.

● COLUMNS   1000

An over-estimate of the number of columns in the constraint matrix
(excluding slacks). Default is 3m where m is specified by the
ROWS card.

● CONJUGATE GRADIENT VERSION   k

This specifies which version of the method of conjugate gradients
should be used if there is not enough storage available (via the
HESSIAN keyword) for the usual quasi-Newton method.

k	Conjugate gradient algorithm
1	Fletcher and Reeves (1964)
2	Polak and Ribiere (1969)
3	Perry (1976)
4	Memoryless DFP
5	Memoryless Complementary DFP

Further discussion is given in Section II.2.

● CRASH OPTION   k

This determines which columns of [A I] are to be considered for

construction of an initial basis B.

k	Meaning
0	The all-slack basis B = I is set up.
1	All columns of A are considered (except those corresponding to nonlinear variables which have been initialized in the MPS data at values away from their bounds -- see Section III.2). This is the default option.
≥ 2	The columns of A corresponding to <u>linear</u> variables will be considered (same as k = 1 if the problem is purely linear).

In all cases the initial basis is chosen to be strictly triangular

without regard to feasibility or optimality.

1. The CRASH option is ignored if a starting basis is specified.

● DEBUG LEVEL   k

Various positive values of k cause various amounts of additional

information to be printed.

● DUMP FILE   k

If k > 0, the last solution obtained will be output to file k in

the format described in Section V.3. Variables are listed by state,

name and value.

● ELEMENTS 3000

An over-estimate of the number of nonzero elements (coefficients, a_{ij}) in the constraint matrix.

1. COEFFICIENTS is a valid alternative keyword.
2. The default value is 5n where n is specified by the COLUMNS card.

● ERROR MESSAGE LIMIT 50

The maximum number of error messages to be printed for each type of error occurring during input. (Should be large for early runs on a particular MPS deck. Can be lowered to suppress warning of non-fatal errors during multiple subsequent runs.)

● FACTORIZATION FREQUENCY k

At most k iterations will occur between factorizations of the basis.

1. INVERT is a valid alternative keyword.
2. Normally k should be around 50 or 60. Values larger than 100 should almost never be used.

● HESSIAN DIMENSION k

This specifies that a $k \times k$ upper-triangular matrix R is to be available for use by the quasi-Newton algorithm, as a means of approximating the reduced Hessian matrix. Suppose that there are ns superbasic variables at a particular iteration.

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1. If $ns \leq k$, the first ns columns of R will be used, although a constant $k(k + 1)/2$ words of storage are allocated from the beginning.

2. The storage for R is substantial if k is as large as 100 or 200. In general, k should be a slight overestimate of the final number of superbasic variables. If the number of nonlinear variables in the objective is nn then k need not be larger than $nn + 1$.

3. If ns exceeds k the quasi-Newton algorithm is discontinued in favor of a conjugate-gradient algorithm. The matrix R is not used until ns decreases to the value k.

4. The default value for k is n, where n is specified by a SUPERBASICS card, or 0 if the problem is linear.

● **IMBED**   YES

   IMBED   NO

Determines whether part of the LU factorization of the basis is to be imbedded in the constraint matrix A.

1. If YES (the default), any non-spike column in B will be referenced in the L file by a single pointer to the appropriate column of A.

2. If NO, any non-spike column will be copied into the L file. This will increase the storage required for L but in general will reduce paging activity on a machine with virtual memory. Execution speed may therefore increase.

● INPUT FILE   k

The file number for the MPS data file. The default value (k = 5)
refers to the card input stream.

● INSERT FILE   k

References a file containing basis information in the format described
in Section V.2.

1. Uses data produced on a PUNCH file.

2. Will be ignored if a positive OLD BASIS FILE is specified.

● INVERT FREQUENCY   k
See FACTORIZATION FREQUENCY

● ITERATIONS   k

An upper limit on the number of iterations allowed.

1. ITNS is a valid alternative keyword.

2. k = 0 is valid; feasibility is checked but optimality is not.

3. The default value is $3m + 10nn$ where m and nn are specified
    by ROWS and NONLINEAR VARIABLES respectively.

● LINESEARCH TOLERANCE   t

For nonlinear problems, this controls the accuracy with which a
minimum of the objective function will be located along the direction
of search each iteration.

1. t must be a real number in the range $0.0 \leq t < 1.0$.

2. Smaller values of t generally result in more evaluations of
    the objective function but fewer total iterations.

3. The default value $t = 0.01$ requests a fairly accurate search, on the assumption that the total work per iteration is large relative to the work required in evaluating the objective function.

4. If the objective is expensive to evaluate, t should be in the range 0.1 to 0.9 (say).

5. The linesearch procedure[†] uses t in the following way. If d is the current search direction and $g(x + \alpha d)$ is the objective gradient at a step α along direction d, then the point $x + \alpha d$ will be accepted as an improvement over the point x if

$$\left| \frac{g(x + \alpha d)^T d}{g(x)^T d} \right| \leq t$$

● LIST LIMIT   100

The maximum number of lines of MPS data to be listed during input. (The header cards, i.e., NAME, ROWS, COLUMNS, RHS, RANGE, BOUNDS, ENDATA, will always be listed along with their position in the MPS deck.)

● LOAD FILE   k

References a file containing basis information in the format described in Section V.3.

1. Uses data produced on a DUMP file.

2. Will be ignored if a positive OLD BASIS FILE or INSERT FILE is specified.

[†]MINOS uses subroutines LNSRCH and NEWPTC from the NPL Algorithms Library. See Gill, Murray, Picken, Barber and Wright (1976).

● LOG FREQUENCY   k

One line of the iteration log will be printed every k-th iteration.
k = 10 or 20 is suggested for those interested only in the final
solution.

● LOWER BOUND   b

The default lower bound to be set on all structural variables before
any bound set is input from the MPS file.

1. May be useful for bounding all variables away from singularities
   in the objective. (Other explicit bounds may also be necessary.)

2. If all or most variables are to be FREE, the negative value $b = -1.0E+30$ (or less) may be used to specify "minus infinity."

● LU ROW TOLERANCE   t

  LU COL TOLERANCE   t

  LU MOD TOLERANCE   t

Tolerances for computing and updating LU factors of the basis. See
Section VIII.

● MAXIMIZE

  MINIMIZE

The required direction of optimization (MIN is default). This
applies to both linear and nonlinear terms in the objective.

● MULTIPLE PRICE   k

If k > 0, this over-rides the PARTIAL PRICE parameter.

1. Whenever a PRICE operation is performed, the k best nonbasic
   variables will be selected for admission to the superbasic set.

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2. <u>Warning</u>: If $k \geq 1$, this puts MINOS into "reduced gradient" mode even on purely linear problems. The subsequent iterations do <u>not</u> correspond to the very efficient "Minor Iterations" carried out by standard LP systems using multiple pricing. (All superbasics will be varied simultaneously. Note however that total storage requirements are independent of k; hence k need not be limited to 5 or 6 as it is in standard LP systems, which require k vectors of dimension m.)

3. On very large nonlinear problems it may be important to use quite a large value of k for the first few hundred or thousand iterations, as a means of "getting a lot of variables going" quickly. For example, if a problem has 3000 variables and 500 of them are nonlinear, the optimal solution may well have 200 or more variables superbasic. If the problem is solved in several runs, it may be beneficial to use $k = 10$ or 20 (say) for early runs, until it seems that the number of superbasics has levelled off.

4. The default strategy (no partial or multiple pricing, add one variable at a time to the superbasics) appears in practice to be preferable once a near-optimal solution has been obtained. This is partly because the projected Hessian approximation for s superbasics is likely to be improved more quickly for an expanded set of $s + k$ superbasics if k is only 1.

5. MULTIPLE PRICE 1 has a slightly different effect from PARTIAL PRICE 1, MULTIPLE PRICE 0, even though just 1 superbasic is added at each PRICE operation. For linear problems, the reduced-gradient mode of operation is slightly less efficient than the simplex algorithm. For nonlinear problems, the method of updating the reduced Hessian approximation differs. With MULTIPLE PRICE 1, the triangular matrix R is expanded to dimension $s + 1$ by appending the <u>unit vector</u> e_{s+1}. (Similarly if $k > 1$, k unit vectors are added.) Otherwise an attempt is made to obtain a better approximation, using some finite-difference technique.

● *NEW BASIS FILE*   k

If $k > 0$, a bit map will be saved on file k every j-th iteration, where j is the specified SAVE FREQUENCY.

1. The first card of the file will contain the word PROCEEDING if the current run is still in progress (see Section V.1).
2. A bit map will also be saved at the end of a run (if $k > 0$).
3. This is the only basis file affected by the SAVE FREQUENCY parameter.

● NONLINEAR VARIABLES   N

This specifies that the first N variables in the constraint matrix are to be regarded as "nonlinear" and made available as the array X(N) to subroutine CALCFG.

1. If $N = 0$, the problem is a linear program.

2. Not <u>all</u> N variables need be explicitly involved in the non-linear objective. For example, some of them may be FIXED at permanent values. Note however that gradient values $G(j)$ must be supplied for all $j = 1, \ldots, N$. The correct value is 0.0 if variable j is essentially a linear variable, or c_j if the linear term $c_j x_j$ is not already included in the linear objective row.

3. In some cases it is convenient to regard even the logical variables as nonlinear (e.g., in solving the linear complementarity problem: $\min z^T w$ subject to $w = Mz + q$, $z^T w = 0$, $z \geq 0$, $w \geq 0$. This problem should be entered as a linear program of the form $Mz \geq -q$, $z \geq 0$, so that the "slack variables" may serve as $-w$.) The value of N should then be $n + 1 + m$ (allowing 1 for the RHS).

● OBJECTIVE   COST

The 8-character name of the free row to be used as the linear part of the objective function.

1. If no OBJECTIVE card exists, or if the name specified is blank, the <u>first</u> free row will be selected.

2. If there are no free rows, or if the user (perhaps intentionally) specifies a non-existent OBJECTIVE name, optimization will terminate at the first feasible solution (unless a nonlinear objective exists).

3. If a nonlinear objective function is of the form $f(x_N) + c^T x$, where $x = (x_N, x_L)$ is partitioned into nonlinear and linear variables respectively, the elements of c may be specified in several alternative ways. Let $c^T x = c_N^T x_N + c_L^T x_L$. Then

    (a) all of c may be specified in the OBJ row of the MPS deck;

    (b) if $c_N^T x_N$ has been programmed in subroutine CALCFG as part of the function $f(x_N)$, then only c_L should be specified in the MPS file;

    (c) either c_N or c_L or both may be absent.

● OLD BASIS FILE   k

Specifies a bit map to be used for the starting solution (see Section V.1).

1. $k > 0$ takes precedence over an INSERT file or a LOAD file.

● PARTIAL PRICE   k

The matrix A will be divided into k equal partitions $A_0, \ldots, A_{k-1}$ for the purposes of selecting a column to enter the basis. Normally one partition A_j will be priced, along with all of the slacks.

1. Almost always k should be set to 1, unless n is large and $m \ll n$, where A is m by n. The default value is $k = 1$.

2. Pricing stops if a variable in A_j or a slack variable has a favorable reduced gradient (as judged by a dynamic tolerance).

3. The next pricing starts with partition $A_{j+1 (\mathrm{mod}\ k)}$.

● PROBLEM NO.   3

For nonlinear problems, this allows the user to set the parameter

"NPROB" in subroutine CALCFG.   See Section III.3.

● PUNCH FILE   k

If k > 0, the final solution obtained will be output to file k

in the format described in Section V.2.   For linear problems, this

format is compatible with various commercial systems.

● RANGES   RANGE001

The 8-character name of the range set to be selected from the MPS file.

1. RNG (or RNGS) are valid alternative keywords.

2. For multiple RANGE sets, see notes 2 and 3 under keyword

    BOUNDS and interpret appropriately.

● REDUCED GRADIENT TOLERANCE   t

For nonlinear problems, this controls the extent to which optimiza-

tion is restricted to the current set of basic and superbasic variables

(i.e., Phase 4 iterations), before a new nonbasic variable is added

to the superbasic set (Phase 3).

1. t must be a real number in the range $0.0 < t \leq 1.0$.   The

    default value is $t = 0.2$.

2. Smaller values of t generally result in more iterations.

    Larger values may be satisfactory, particularly when the current

    solution is far from optimal.

3. Note: This parameter does not refer to the absolute size of the reduced gradient at any stage. It is not a means for requesting iterations to terminate when the reduced gradient reaches a certain value.

4. t is used in the following way. Whenever a variable x_j is added to the superbasic set the norm of the reduced-gradient vector (for all superbasics) is recorded. Let this be $\| \hat{g}_0 \|$. Usually $\| \hat{g}_0 \| = |d_j|$, where d_j is the "reduced cost" for the new superbasic. Subsequent Phase 4 iterations will continue at least until the norm of the reduced-gradient vector satisfies $\| \hat{g} \| \leq t \| \hat{g}_0 \|$. (Other convergence tests may cause further Phase 4 iterations.)

● RHS   RHSIDE1

The 8-character name of the rhs vector to be selected from the MPS data.

1. See notes 2 and 3 under BOUNDS.

2. If no rhs exists after input (perhaps after deletion of small coefficients), a dummy zero vector is constructed.

● ROWS   500 .

An over-estimate of the number of rows in the constraint matrix.

1. If m is the number specified, the size of the hash table used to hold the row names will usually be $2m$. The number of collisions (and hence total input time) may in general be reduced by specifying m to be larger than necessary (within reason).

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● SAVE FREQUENCY   k

Suppose the NEW BASIS FILE card specified a file number f. If
$f > 0$ and $k > 0$, a bit map describing the current solution will
be saved on file f every k-th iteration.

● SOLUTION   w

Determines whether a solution is to be printed on file 6 and/or
output to a specified file, as follows:

w	Meaning
YES	The last solution obtained will be printed on file 6 (whether optimal or not).
NO	The solution will not be printed.
IF OPTIMAL	The solution will be printed only if it has been determined to be optimal, or if the problem appears to be infeasible or unbounded.
FILE k	If $k > 0$ the last solution obtained will be output to file k (whether optimal or not).

1. If w is blank or is anything other than YES, NO or FILE it
   will be treated as IF OPTIMAL.

2. The first 3 options (YES, NO, IF OPTIMAL) may co-exist with
   the FILE option.

3. For the 1st and 3rd options, floating-point numbers are printed
   with F16.5 format, and "infinite" upper or lower bounds are
   denoted by the word NONE.

4. For the FILE option, the format used is 1PE16.6 for all numbers, including "infinite" bounds which will appear as ±1.000000E+30.

5. Since the printer is file 6, the specifications

SOLUTION   NO

SOLUTION FILE   6

may be used together to obtain E-format in the printed solution (e.g., if 7 significant digits are required for all numbers).

● SUPERBASICS   ns

Specifies how many superbasic variables are to be allowed for. Should be used in conjunction with a "HESSIAN DIMENSION" card.

1. In some sense, ns measures the expected deviation from linearity.

2. Normally ns need not be greater than nn + 1 where nn is the specified number of nonlinear variables.

3. For certain problems, ns may be considerably smaller than nn. This is important if nn is very large.

● TARGET OBJECTIVE VALUE   t

This keyword invokes a certain optimization technique, and also sets the value of t. In effect it forces the solution at each iteration to be regarded as infeasible if the linear objective row has not yet reached the value t. This is designed to take advantage of the

special pivot-row selection procedure (CHUZR) due to D.C. Rarick (1975), which minimizes the _sum_ of infeasibilities at each iteration without concerning itself with the _number_ of infeasibilities.

1. TARGET has no effect if there is no linear objective.

2. The value of t should take into account the direction of optimization (MIN or MAX).

3. It is generally unwise to invoke targeting unless t is known a priori to be sub-optimal.

4. If the problem appears to be infeasible, targeting is disabled and optimization continues normally.

5. Stay clear unless you know what you are doing.

● UPPER BOUND   b

The default upper bound to be set on all structural variables before any bound set is input from the MPS file.

1. A value of 1.0E+30 or greater will be interpreted as "plus infinity." This is the default value and therefore need not be set.

● VERIFY   w

Determines whether or not the computed gradient vector from the user's subroutine CALCFG is to be checked numerically. If w is the word NO, the gradient will not be checked. Any other value of w is equivalent to YES. Checking will be performed at the first feasible point found. Finite differences are taken of the objective value along two orthogonal directions. (If necessary, the directions will be modified to avoid stepping outside the feasible region.) The default value of w is NO.

● WEIGHT ON OBJECTIVE   w

This keyword invokes the so-called <u>composite</u> <u>objective</u> technique,
if the first solution obtained is infeasible. The method attempts
to optimize the true objective function (linear terms only) while
simultaneously reducing the sum of infeasibilities.

1. The objective function is defined at each infeasible iteration
    to be

$$\min \, mw(c^T x) + (\text{sum of infeasibilities})$$

    where m = 1 for MIN, -1 for MAX, and c is the linear
    objective row.

2. If the sum of infeasibilities can not be reduced at some stage,
    w is reduced by a factor of 10. This helps to allow for the
    possibility that the w specified by the user is too large.
    It also provides dynamic allowance for the fact that the sum of
    infeasibilities is tending towards zero.

3. The effect of w is disabled after 5 such reductions of w,
    or if a feasible solution is obtained.

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V. Basis File Formats

Figures 1 through 4 illustrate the data formats used for the various basis files introduced in Section III.4. Selected column numbers are included to define the significant data fields. The logical records in each file are alphanumeric data with the following record lengths:

File	Logical record length
OLD/NEW BIT MAP	80
PUNCH/INSERT FILE	36
DUMP/LOAD FILE	36
SOLUTION FILE	108

(Larger record lengths may be used for convenience, e.g., 80, 80, 80, 133 respectively.)

The files shown correspond to the optimal solution for a small but non-trivial example, namely the Weapon Assignment problem in Himmelblau (1972), problem 23. This is of the form

$$\text{minimize} \quad \sum_{j=1}^{20} u_j \left(\prod_{i=1}^{5} a_{ij}^{x_{ij}} - 1 \right)$$

$$\text{subject to} \quad Ax \gtrless b, \qquad x \geq 0$$

where A has dimensions 12×100.

V.1. OLD/NEW Basis Files (bit maps)

A bit map is primarily intended for restarting the solution of a problem at a point that was reached by an earlier run, either on the _same problem_ or on a related problem of the _same dimensions,_ e.g., one which has a different OBJECTIVE, RHS, RANGE or BOUND set or different matrix coefficients.

```
WEAPON A    ITN   122      OPTIMAL SOLN   PHASE  3        OBJ  -1.7356957785560+03
OBJ=         RHS=B          RNG=          BMD=            M=  12  N= 113  SB= 18
000032200020000200002000020000300002000020000200002000020000200002000020000200002030
0220000200002000020000300002000010000003122123
    41  2.0316768382961900+01
    82  3.3027533814957700+00
    17  2.3489533295527300+01
    36  2.7066733129639200+01
    59  4.0934258239765500+01
    22  2.0899605737884000+01
    15  4.8628858754640800+01
    88  5.7551539944999450+01
    65  5.4015197600566400+01
   112 -6.2810931659367300+00
    93  6.4211504061437200+01
   109 -1.1132295900964500+01
     6  1.3516456395590100+01
    72  2.4230467215620100+01
   110 -8.8237529151783900+00
    87  7.2033835250449810+01
     7  1.3604576472287700+00
    10  4.5408197027647600+01
     0
```

Figure 1. Format of OLD and NEW BASIS FILES (Exit Data)

The information saved is very compact (but rather difficult to modify). With some computer systems it is possible to access the bit map at occasional intervals during a run; in such cases the first card of the file will provide a useful glimpse of the run's progress so far.

As illustrated in Figure 1, the information contained in an *OLD* or *NEW BASIS FILE* is as follows:

1. A card containing the problem name, iteration number, current status (OPTIMAL SOLN, INFEASIBLE, UNBOUNDED, EXCESS ITNS, ERROR CONDN, or PROCEEDING), current phase and current objective value or sum of infeasibilities.

2. A card containing the OBJ, RHS, RNG and BND names, M = no. of rows in the constraint matrix, N = total no. of variables, SB = current no. of superbasics. In this context,

$$N = n + 1 + m$$

where

(a) the first n variables are those in A;

(b) variable $n + 1$ is the RHS (fixed at -1.0);

(c) the last m variables are the logicals (slacks).

3. A set of $(N-1)/80 + 1$ cards indicating the state of each of the above N variables. One character $HS(j)$ is recorded for each $j = 1, 2, \ldots, N$ as follows, written with FORMAT(80I1):

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HS(j)	State of the j-th variable
0	Nonbasic at lower bound
1	Nonbasic at upper bound
2	Superbasic
3	Basic

If variable j is FIXED (lower bound = upper bound), then $HS(j)$ may be 0 or 1. The same is true if variable j is FREE (infinite bounds) and still nonbasic, although free variables will usually be basic.

4. A set of cards of the form

$$j \qquad x_j$$

where j is a variable number and x_j is a value. These are written with FORMAT(I8, 1PE24.14) and are terminated by a card with $j = 0$. Normally j will be in the range 1 to N and will correspond to a variable for which $HS(j) = 2$ (superbasic).

Notes on Input of an OLD BASIS FILE

1. Card 1 above is printed but not otherwise used.

2. The values labelled M and N on card 2 must agree with those for the MPS file that has just been read. The value labelled SB is read and printed as an integer but is not otherwise used.

3. Cards 3, 4, etc., must contain exactly M values $HS(j) = 3$ (for the basic variables).

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4. The cards containing j and x_j values need not occur in natural column order (e.g., see Figure 1). The user may modify or extend this list of cards relatively easily as long as the following points are borne in mind.

5. The list should include a card for all variables whose state is $HS(j) = 2$ (superbasic).

6. For any given j and x_j, if $HS(j) \neq 3$ then variable j will normally be initialized at the value x_j and its state will be reset to $HS(j) = 2$ (superbasic).

7. If $HS(j) = 3$, the value x_j will be recorded for nonlinear variables but no other action will be taken (thus variable j will remain basic).

8. If the specified value x_j lies on or outside the bounds on variable j, or if the total number of superbasics has already reached the limit specified in the SPECS file, then variable j will be made nonbasic at the bound <u>nearest</u> to x_j.

V.2. <u>PUNCH and INSERT files (industry standard)</u>

These files specify a basis by a list of cards of the form

INDICATOR   NAME1   NAME2   VALUE

as shown in Figure 2. The output to a PUNCH file may be used directly as input from an INSERT file, either for MINOS or (if the problem is purely linear) for one of the commercial LP systems. The various indicators are best defined in terms of the action taken on input. It is assumed that the basis is initially the set of logical variables and that the structurals are nonbasic at their smallest bound in absolute magnitude.

```
NAME             WEAPON A   PUNCH/INSERT
 XL X015         L1          5.081550+01
 SB X021                     1.351650+01
 SB X022                     1.360460+00
 SB X025                     4.540820+01
 SB X035                     4.862890+01
 SB X042                     2.348950+01
 SB X052                     2.089960+01
 XL X061         L2          1.000000+02
 XL X071         L3          3.910000+01
 SB X081                     2.706670+01
 SB X091                     2.031680+01
 XL X105         L4          5.113230+01
 XL X114         L5          3.319740+01
 SB X124                     4.093430+01
 SB X135                     5.401520+01
 XU X144         G7          5.882380+01
 XL X152         G8          2.621120+01
 XL X153         G9          4.378880+01
 SB X162                     2.423650+01
 XU X164         G10         1.704460+01
 SB X172                     3.802750+00
 SB X173                     7.203380+01
 SB X183                     5.755150+01
 SB X193                     6.421150+01
 XL X203         G11         6.241430+01
 SB G8                      -1.113230+01
 SB G9                      -8.823750+00
 SB G11                     -6.281090+00
ENDATA
```

```
  |    |      |   |       |  |            |
  1    5      12  15      22 25           36
```

Figure 2.   Format of PUNCH and INSERT FILES

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INDICATOR	Action Taken During INSERT
XL (XU)	Make structural NAME1 basic and logical NAME2 nonbasic at its lower (upper) bound.
LL (UL)	Make variable NAME1 nonbasic at its lower (upper) bound.
SB	Make variable NAME1 superbasic at the value VALUE.

Notes on PUNCH Data

1. Variables are output in natural order. For example, on the first XL or XU card, NAME1 will be the first basic structural and NAME2 will be the first logical which is not basic (it may be nonbasic or superbasic).

2. Indicator SB is an addition to the standard format to allow for non-vertex solutions.

3. Superbasic logicals are output last. These will be some of the logicals which previously appeared as NAME2 (e.g., logicals G8, G9, G11 in Figure 2).

4. PUNCH and INSERT files deal with the status and value of logical variables, whereas SOLUTION files deal with row status and activities.

5. LL or UL cards are not output for nonbasic variables if the corresponding bound value is zero.

Notes on INSERT Data

1. Preferably an INSERT file should be an unmodified PUNCH file obtained from an earlier run on the same problem.

2. VALUE is used only if INDICATOR is SB or NAME1 is a nonlinear variable.

3. _Warning_ -- if D or E format is used for a VALUE, the exponent must be right-justified to column 36.

4. Additional SB cards may be added before the ENDATA card, to specify additional superbasic structurals or logicals. (Any such card will be ignored if previous cards have implied that NAME1 is basic or superbasic, or that NAME2 is not basic.)

5. If the specified VALUE on an SB card lies on or outside the bounds on variable NAME1, or if the total number of superbasics has already reached the limit specified in the SPECS files, then variable NAME1 will be made nonbasic at the bound _nearest_ to VALUE.

V.3. DUMP and LOAD Files

These files are similar to PUNCH and INSERT files but record solution information in a manner that is more direct and more easily modified. In particular, no distinction is made between row and column variables. (Note that status indicators and values refer to _rows_ as in the SOLUTION file, rather than the corresponding logical variables.)

```
NAME                 WEAPON A         DUMP/LOAD
  BS  X015                          5.081550+01
  SB  X021                          1.351650+01
  SB  X022                          1.360460+00
  SB  X025                          4.540820+01
  SB  X035                          4.862890+01
  SB  X042                          2.348950+01
  SB  X052                          2.089960+01
  BS  X061                          1.000000+02
  BS  X071                          3.910000+01
  SB  X081                          2.706670+01
  SB  X091                          2.031680+01
  BS  X105                          5.113230+01
  BS  X114                          3.319740+01
  SB  X124                          4.093430+01
  SB  X135                          5.401520+01
  BS  X144                          5.882380+01
  BS  X152                          2.621120+01
  BS  X153                          4.378380+01
  SB  X162                          2.423650+01
  BS  X164                          1.704460+01
  SB  X172                          3.802750+00
  SB  X173                          7.203380+01
  SB  X183                          5.755150+01
  SB  X193                          6.421150+01
  BS  X203                          6.241430+01
  BS  G6                          -2.081550+01
  SB  G8                          -1.113230+01
  SB  G9                          -8.823750+00
  SB  G11                         -6.281090+00
  BS  G12                         -5.241430+01
ENDATA
```

```
:     :        :              :            :
1     5        12             25           36
```

Figure 3.   Format of DUMP and LOAD FILEs

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Each data card is of the form

INDICATOR   NAME   (blank)   VALUE

as shown in Figure 3. The indicators LL, UL, BS and SB mean
Lower Limit, Upper Limit, Basic and Superbasic respectively.

Notes on DUMP Data

1. Variables are output in natural order, columns first, then rows.

2. LL or UL cards are not output for nonbasic variables if the
    corresponding bound value is zero, or if the variable NAME is
    FIXED (lower bound = upper bound).

3. Nonbasic FREE variables will be output with either LL or
    UL indicators but with value zero.

Notes on LOAD Data

1. Normally a DUMP file will be valid as a LOAD file to
    restart solution of the same problem.

2. VALUE is used only if INDICATOR is SB or if NAME is a
    nonlinear variable.

3. If the VALUE on an SB card lies on or outside the bounds
    on variable NAME, or if the total number of superbasics has
    already reached the limit specified in the SPECS file, then
    variable NAME will be made nonbasic at the bound nearest to
    VALUE.

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4. (Partial basis) If fewer than M variables are specified to be basic, a tentative basis list will be constructed by adding the requisite number of logicals, starting from the first row and taking those that were not previously specified to be basic or superbasic.

Note: If the resulting basis is singular, the basis factorization routine will reject a number of variables and introduce an appropriate set of logicals. The rejected variables will be favored during subsequent iterations for readmission to the basis, but otherwise the starting point obtained will not necessarily be "good."

5. (Too many basics) If M variables have already been specified as basic, any further BS indicators will be treated as though they were SB. This feature may be useful in combining solutions to smaller subproblems.

V.4. SOLUTION File

Figure 4 illustrates the format of the file output at the end of a run if requested by a SPECS card of the form

SOLUTION FILE   k

$(k > 0)$.

PROBLEM NAME: WEAPON A   OBJECTIVE VALUE: -1.735369579D+03

STATUS: OPTIMAL SOLN   PHASE: 3   ITERATION: 122

OBJECTIVE: 0 (MIN)
RHS:
RANGES:
BOUNDS:

SECTION 1 - ROWS

NUMBER	...ROW..	AT	...ACTIVITY...	SLACK ACTIVITY	..LOWER LIMIT.	..UPPER LIMIT.	.DUAL ACTIVITY	..I
102	L1	UL	2.000000D+02	0.0	-1.000000E+30	2.000000E+02	-5.992750D-02	1
103	L2	UL	1.000000D+02	0.0	-1.000000E+30	1.000000E+02	-2.176944D-01	2
104	L3	UL	3.000000D+02	0.0	-1.000000E+30	3.000000E+02	-6.870748D-02	3
105	L4	UL	1.500000D+02	0.0	-1.000000E+30	1.500000E+02	-1.235846D-01	4
106	L5	UL	2.500000D+02	0.0	-1.000000E+30	2.500000E+02	-7.299930D-02	5
107	L6	BS	5.081545D+01	-2.081545D+01	3.000000E+01	1.000000E+30	0.0	6
108	G7	LL	1.000000D+02	0.0	1.000000E+02	1.000000E+30	5.992664D-02	7
109	G8	SBS	5.113230D+01	-1.113230D+01	4.000000E+01	1.000000E+30	7.836420D-09	8
110	G9	SBS	7.023750D+01	-8.623750D+00	5.000000E+01	1.000000E+30	-1.209550D-09	9
111	G10	LL	7.000000D+01	0.0	7.000000E+01	1.000000E+30	-2.098790D-02	10
112	G11	SBS	4.120109D+01	-6.281093D+00	3.000000E+01	1.000000E+30	-1.884939D-09	11
113	G12	BS	6.241430D+01	-5.241430D+01	1.000000E+01	1.000000E+30	0.0	12

Figure 4. Format of the SOLUTION File

SECTION 2 - COLUMNS

NUMBER	.COLUMN.	AT	...ACTIVITY...	..OBJ GRADIENT.	..LOWER LIMIT.	..UPPER LIMIT.	.REDUCED COST.	M+J
1	X011	LL	0.0	0.0	0.0	1.00000E+30	5.992750D-02	13
2	X012	LL	0.0	-1.511620D-01	0.0	1.00000E+30	6.553250D-02	14
3	X013	LL	0.0	-3.039220D-02	0.0	1.00000E+30	3.331527D-02	15
4	X015	LL	5.301545D+01	-7.229090D-02	0.0	1.00000E+30	1.235846D-01	16
5	X021	BS	1.351846D+01	-5.992750D-02	0.0	1.00000E+30	0.0	17
6	X022	BS	1.300490D+00	-2.175944D-01	0.0	1.00000E+30	0.0	18
7	X023		0.0	-5.992750D-02	0.0	1.00000E+30	8.779980D-03	19
8	X024		0.0	0.0	0.0	1.00000E+30	1.258460D-01	20
9	X025	BS	4.044020D+01	-7.229090D-02	0.0	1.00000E+30	5.992750D-02	21
10	X031	LL	0.0	-1.409021D-01	0.0	1.00000E+30	7.572390D-02	22
11	X032		0.0	-3.039220D-02	0.0	1.00000E+30	3.331527D-02	23
12	X033		0.0	-7.229090D-02	0.0	1.00000E+30	1.235846D-01	24
13	X034		4.862862D+01		0.0	1.00000E+30	0.0	25
14	X035		2.040752D+01	-2.176945D-01	0.0	1.00000E+30	5.992750D-02	26
15	X041		0.0	-5.992950D-02	0.0	1.00000E+30	1.773787D-02	27
16	X042		0.0	0.0	0.0	1.00000E+30	1.235846D-01	28
17	X043		2.009010D+01	-2.176944D-01	0.0	1.00000E+30	2.472740D-03	29
18	X044		0.0	-5.468114D-02	0.0	1.00000E+30	5.992750D-02	30
19	X045		1.003555D+02	-0.607470D-02	0.0	1.00000E+30	0.0	31
20	X051		0.0	-8.529750D-07	0.0	1.00000E+30	1.406330D-02	32
21	X052		0.0	-1.105910D-06	0.0	1.00000E+30	1.235846D-01	33
22	X053		0.0	-5.969950D-07	0.0	1.00000E+30	3.583430D-03	34
23	X055		3.910043D+01	-1.003030D-07	0.0	1.00000E+30	1.575570D-01	35
24	X061		0.0	-5.992750D-02	0.0	1.00000E+30	8.780620D-03	36
25	X062		0.0	-1.198550D-01	0.0	1.00000E+30	0.355790D-02	37
26	X063		0.0	-4.742022D-02	0.0	1.00000E+30	1.264150D-02	38
27	X064		6.766075D+01	-1.149100D-02	0.0	1.00000E+30	9.783945D-02	39
28	X071		0.0	-5.992743D-02	0.0	1.00000E+30	2.181260D-02	40
29	X072		0.0	-3.477627D-02	0.0	1.00000E+30	1.235846D-01	41
30	X074		5.816077D+01	0.0	0.0	1.00000E+30	0.079990D-02	42
31	X075		0.0	-5.992743D-02	0.0	1.00000E+30	0.0	43
32	X081		0.0	-5.992743D-02	0.0	1.00000E+30	1.451730D-01	44
33	X082		0.0	-2.392990D-02	0.0	1.00000E+30	3.393120D-01	45
34	X083		0.0	0.0	0.0	1.00000E+30	1.235846D-01	46
35	X084		5.814350D+01	-1.604830D-01	0.0	1.00000E+30	7.229090D-02	47
36	X085		0.0	-3.519381D-02	0.0	1.00000E+30	1.421050D-01	48
37	X091		0.0	-2.800745D-02	0.0	1.00000E+30	1.770700D-01	49
38	X092		0.0	-7.229087D-02	0.0	1.00000E+30	4.314490D-02	50
39	X093		2.317350D+01	0.0	0.0	1.00000E+30	7.229090D-02	51
40	X095		0.0	-1.316755D-02	0.0	1.00000E+30	5.227490D-01	52
41	X101		0.0	-1.235845D-01	0.0	1.00000E+30	3.351306D-02	53
42	X102		0.0	0.721407D-02	0.0	1.00000E+30	9.557544D-02	54

The format of a _printed_ solution is essentially identical except for the following:

1. The first character in each line is for carriage control and will not appear on the printed page. Only two lines have a non-blank carriage control.

2. Format F16.5 is used for numerical values rather than 1PE16.6.

3. Any UPPER or LOWER LIMITS that are "infinite" ($\pm$ 1.0E+30) appear as the word NONE.

In the ROWS section, both ROW and SLACK activities are shown. The quantities labelled DUAL ACTIVITY are the _simplex multipliers_, π.

In the COLUMNS section, OBJ GRADIENT means, for linear variables, the coefficient in the objective row, say c_j. For nonlinear variables, if the solution is feasible, it means $g_j + c_j$ where g_j is the relevant component of the gradient of the non-linear function $f(x)$. (If the solution is infeasible, just c_j appears.)

The quantities labelled REDUCED COST are the reduced gradients $d_j = c_j - \pi^T a_j$ (or $g_j + c_j - \pi^T a_j$) for each column a_j. At an optimal solution these should be zero for basic variables (status BS), non-negative for variables at LL, non-positive for variables at UL, and zero or very small (in absolute magnitude) for variables labelled SBS.

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If the solution is infeasible, the status of basic and superbasic variables (rows or columns) may be shown as follows:

++   variable is above its upper limit

--   variable is below its lower limit.

A SOLUTION file is intended to be read from disk by a self-contained program which extracts and saves certain values as required, possibly for further computation. Typically the first 14 lines would be skipped. Each subsequent line should be read using a format of the form

FORMAT(I8, 2X, 2A4, A4, 5E16.6, I6)

The end of the ROWS section can be detected by the almost blank line. If this and the next 4 lines are skipped, the COLUMNS section can then be read under the same format. Hence the inefficiency of BACKSPACE or re-read operations can be avoided.

VI. Restarting Modified Problems

The preceding sections document three distinct starting methods (OLD BIT MAP, INSERT, LOAD) which may be preferable to any of the three cold start (CRASH) options. The best choice depends on the extent to which the present problem has been modified, and whether variables are to be specified by name or by number.

1. (Protection) In general the user is protected against accidental specification of infinite bounds. For example, if UL is specified for some variable that has an upper bound of plus infinity, the indicator is treated as if it were LL. Similarly for LL cards, if the variable's lower bound is minus infinity. No message is issued. (Note -- switching bounds is safe even with nonbasic FREE variables. They will be treated as having value zero.)

2. (Default status) If the status of a variable is not explicitly specified, it will initially be nonbasic at the bound which is <u>smallest in absolute magnitude</u>.

3. (Relaxing of bounds) Suppose that in Problem A a certain variable X has bounds

$$0 \leq X \leq 5.0$$

and that when a solution is saved, X is nonbasic at its upper bound 5.0. Certain care must be taken if the user wishes to use the solution to start some Problem B in which the upper

bound on X has been relaxed. If <u>no</u> action is taken, two cases arise:

(a) If the upper bound has been <u>removed</u>, a bound switch would normally occur (e.g., UL would be treated as LL) and X would be made nonbasic at its <u>lower</u> bound, 0.0.

(b) If the bound has been <u>relaxed</u> to some larger but finite value, e.g., $0 \leq X \leq 10.0$, X would normally be made non-basic at that upper bound.

In both cases the new solution will clearly be different from the old. In some situations this may not be of consequence, but often the new solution may prove to be severely infeasible or seriously removed from optimality.

The same comments apply if X were FIXED at 5.0 in Problem A.

To retain the original solution precisely (all other things being equal), the user need only request that X be made super-basic at the required value 5.0. This means inserting a card of the form

$$j \qquad 5.0$$

near the end of a bit map, or an

$$SB \qquad X \qquad 5.0$$

card near the end of an INSERT or LOAD file.

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In all cases, the necessary storage must be requested by a SPECS card of the form

$$\text{SUPERBASICS} \qquad k$$

where k is at least as large as the number of superbasics being specified.

4. In case (a) above the user may happen to know that the upper bound 5.0 in Problem A will effectively be a lower bound on the optimal value of X in Problem B. Of course it would then be simple to define the bounds on X to be

$$5.0 \leq X \leq \infty$$

in the MPS data for Problem B, and no further action need be taken. In case (b), the bounds

$$5.0 \leq X \leq 10.0$$

could be included, and the state of X would have to be changed from UL to LL. However, when many interacting variables are involved it may be difficult to determine in advance that X will move above its old bound; the previous method (making X superbasic) is then much to be preferred.

5. (Tightening of bounds)   If Problems A and B have the bounds

$$0 \leq X \leq 5.0$$

$$0 \leq X \leq 2.0$$

respectively, and if X is again at its upper limit in
Problem A, there is no way in general of initializing X
at value 5.0 for Problem B. MINOS will move X to its
nearest bound, in this case 2.0, and hopefully most users
will be happy. However, note the following rule of thumb.

6. (Other modifications) Wherever possible, a series of problems
should be ordered so that the most tightly constrained cases
are solved _first_. Their solutions will often provide feasible
starting points for subsequent relaxed problems.

VII. The Iteration Log

One line of information is printed on file 6 every k-th iteration, where k is specified by a card of the form LOG FREQUENCY k. The default value is k = 1. For large problems, k = 10 will save a significant volume of paper.

In the following description, a PRICE operation is defined to be the process by which one or more nonbasic variables are selected to become *superbasic*. Normally just one variable is selected, which we will denote by J. If the problem is purely linear, variable J will usually become basic immediately (unless it should happen to reach its opposite bound and return to the nonbasic set).

If Partial Pricing is in effect, variable J is selected from the matrix $[A_{PP} \, I]$ where A_{PP} is the PP-th partition of A, and I corresponds to the slack variables.

If Multiple Pricing is in effect, several variables may be selected from the whole matrix [A I]. In this case J refers to the variable with the largest favorable "reduced cost."

The iteration log includes the following items.

ITN   The current iteration number.

PH   The current phase, as follows:

1 => Phase 1 simplex procedure (the current solution is infeasible).

2 => Phase 2 simplex procedure (e.g., if the problem is a linear program, or if none of the nonlinear variables is basic).

3 => A PRICE operation has just been performed and the current iteration is a non-simplex step.

4 => Optimization is being performed amongst the current basic and superbasic variables (ignoring the nonbasics).

PP   The "Partial Price" indicator, reset by the PRICE operation in Phase 1, 2, or 3.

PP = 1 if partial pricing is not in effect.

PP = 0 if variable J is one of those rejected from the basis by the previous basis factorization.

NOPT   The number of "non-optimal" variables present in whatever set of nonbasics was scanned during the last PRICE operation. Reset in Phase 1, 2, or 3.

DJ/RG   In Phase 1, 2, or 3 this is "DJ," the reduced cost of the variable J selected by PRICE. In Phase 4 it is "RG," the norm of the reduced gradient vector; i.e., the largest reduced gradient amongst the current set of superbasics.

+SBS   The variable J selected by PRICE to become superbasic.

−SBS   The variable chosen to leave the set of superbasics (to become either basic or nonbasic).

−BS   The variable chosen to leave the basis to become nonbasic.

STEP   The distance moved along the current search direction.

PIVOT   If column a_q replaces the p-th column of the basis B, this is the p-th element of the transformed column $B^{-1}a_q$.

NSPK   The current number of spikes in the LU factorization of B. This is the number of non-trivial columns in U (excluding ones that have been deleted in earlier iterations).

L   The number of nonzeros in the L part of the factorization of B (excluding elements that are imbedded in the constraint matrix A).

U   The number of nonzeros in the U part of the factorization of B (including ones in spikes that have been deleted in earlier iterations). This does not include nonzeros in a certain matrix F, which is really part of U but is stored separately as an upper triangular matrix of dimension NSPK.

NINF   The number of infeasibilities before the present iteration. (Note -- this number need not decrease monotonically.)

SINF/   If NINF > 0, this is SINF, the sum of infeasibilities
OBJECTIVE   before the present iteration; it should decrease mono-
              tonically.

              If NINF = 0, the value of the true objective _after_ the
              present iteration is given (including both linear and
              nonlinear terms).

       The following items are printed if the problem is nonlinear
or if the Phase is 3 or 4.

NFG   The number of times the nonlinear objective and its
              gradient have been evaluated.

NSB   The current number of superbasic variables.

R1M   An indication of the type of modifications that have
              been made to the matrix R being used to approximate
              the projected Hessian matrix.

H-CONDN   This is an estimate of the condition number of the
              reduced Hessian matrix. More precisely it is K,
              the square of the ratio of the largest and smallest
              diagonals of the triangular matrix R, which constitutes
              a lower bound on the condition number of the matrix
              $(R^T R)$ that approximates the projected Hessian.
              The value of K gives a rough indication of whether
              or not the optimization procedure is having difficulty.
              Let ϵ be the relative precision of the floating-point
              arithmetic being used. The reduced-gradient algorithm
              will make slow progress if K becomes as large as $\epsilon^{-1/2}$,

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and will probably fail to find a better solution if K reaches $\epsilon^{-3/4}$ or larger. In such cases the user should

(1) check that the constraint matrix is well scaled;

(2) scale the variables involved in the nonlinear objective so that the Hessian matrix is as well conditioned as possible (ignoring rows and columns corresponding to nonbasic variables);

(3) add upper or lower bounds to certain variables to keep them a reasonable distance from any singularities in the objective or its derivatives.

CONV   A set of four logical variables C_1, C_2, C_3, C_4 that are used to determine when to discontinue optimization in the current subspace (Phase 4) and consider moving off a constraint (Phase 3).

Let RG be the current norm of the reduced gradient. The meaning of the C_j is briefly as follows:

C_1 is TRUE if the change in x was sufficiently small.

C_2 is TRUE if the change in $f(x)$ was sufficiently small.

C_3 is TRUE if RG is smaller than some loose tolerance.

C_4 is TRUE if RG is smaller than some tight tolerance.

The test used is of the form _if_ $(C_1$ _and_ C_2 _and_ $C_3)$ _or_ C_4 _then go to_ Phase 3. All tolerances are varied dynamically.

VIII. Tolerances for Computing the Basis Factorization

Three tolerances T_r, T_c and T_m are involved in computing and updating the LU factorization of the basis. They may be set by specification cards of the following form (their default values are shown as comments):

LU ROW TOLERANCE = 0.01   * DEFAULT T_r = 0.001

LU COL TOLERANCE = 0.9   * DEFAULT T_c = 0.1

LU MOD TOLERANCE = 1.0   * DEFAULT T_m = 0.99

The use of T_r and T_c must be understood in terms of the "bump and spike" structure of a basis matrix as introduced by Hellerman and Rarick (1971, 1972) in their Preassigned Pivot Procedures P^3 and P^4. The general aim is to accept the preassigned pivot order whenever possible (since it has close to the minimal number of spikes) but to revise this order where necessary to maintain numerical stability.

In this context, attention is restricted to the rows and columns inside any one bump, and the diagonal elements between bumps are regarded as bumps of dimension 1. Each column in a bump is initially either a _spike_ column or a _triangle_ (non-spike)column.

The k-th column in a bump fails a _row_ test if its preassigned pivot element is smaller than T_r times the largest remaining potential pivot in that column. Some replacement column is then selected from the remaining spikes in the bump as follows. The

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pivot elements α_j are computed for each remaining spike and the largest, α_{max}, is determined. Let column s be the <u>first</u> spike such that α_s is larger than $T_c * \alpha_{max}$. Then columns k and s are interchanged.[†]

The above process amounts to Gaussian elimination with column interchanges.

Increasing T_r towards 1.0 will tend to

(a) increase the amount of column interchanging (spike-swapping);

(b) increase the total no. of spikes (if column k above is a triangle column the spike count will increase by 1);

(c) increase the density of the LU factorization (fill-in occurs only in spike columns);

(c) improve numerical stability.

Increasing T_c towards 1.0 will tend to

(a) increase (slightly) the density of the LU factors (early spikes are generally shorter and will have less fill-in than later spikes),

(b) improve numerical stability.

[†] If there is no suitable substitute, column k will be replaced by an appropriate slack vector.

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For large-scale problems the total number of spikes is critical and a compromise must be reached between sparsity and stability. The default values for T_r and T_c are safe for most problems. If numerical difficulties arise, T_r can be increased to 0.01 and T_c to 0.9 or 0.99 without serious loss of efficiency. Larger values for T_r may degrade performance significantly.

For small dense problems (i.e., "non-LP" problems such as those arising from data-fitting), maximum stability is readily obtained by setting both T_r and T_c to 0.99 (say). The loss of efficiency will be negligible.

Stability during _updating_ of the LU factorization is controlled by the tolerance T_m. In this implementation of the method of Bartels and Golub, maximum stability is obtained via the default value $T_m = 0.99$ with essentially _no loss of efficiency_. There is normally no reason to use a smaller tolerance. However, one can simulate the numerical properties of the method of Forrest and Tomlin (1972) by setting $T_m = 0.0$.

IX.   System Information -- Machine-Dependent Matters

The source code for MINOS is approximately 8300 cards of standard Fortran IV. One of the design goals has been to maintain compatibility with the Fortran compilers on the principal large scientific machines (e.g., IBM, CDC, Univac, Burroughs). At the same time it has been possible to maintain a reasonably high level of efficiency and flexibility. For example, problems of arbitrary size may be solved without recompilation of the program, and the method for allocating storage can take maximum advantage of the available word sizes on a particular machine.

The main items involved in installing the system on a machine other than the IBM Systems 360 and 370 are:

--The setting of 4 variables in subroutine INITIZ.

--The hash function used in subroutine HASH.

--The file numbers used for the reader, printer and scratch files.

These topics are documented in the following sections.

IX.1.   Precision, Word Sizes

Four scalar quantities must be initialized in subroutine INITIZ to match any particular machine. Their values depend on whether single or double precision floating-point arithmetic is required throughout, and their meaning is best understood in terms of the following values which are appropriate for the IBM System/370:

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$$EPS = 16.0^{**}(-13)$$
$$NWORDR = 2$$
$$NWORDI = 2$$
$$NWORDH = 4$$

The variables in all relevant subroutines are implicitly typed as follows:

```
IMPLICIT    REAL(A-B), REAL*8(C-G, O-Z), INTEGER(I-N), INTEGER*2(H)
```

Furthermore, all array storage is provided by segmenting a single array of the form

```
REAL*8  Z(NWCORE)
```

For the IBM 370, the four critical values above are therefore interpreted as follows:

1. EPS contains the smallest positive number such that

$$W = 1.0 + EPS$$

gives W a value greater than 1.0. (We assume that the Fortran compiler will convert 1.0 to double precision and perform the addition in double precision.)

2. NWORDR = 2 means that if some array is of type REAL then 2 elements of that array occupy the same storage as one element of the REAL*8 array Z. Similarly, for arrays of type INTEGER and INTEGER*2, the respective numbers are 2 and 4 elements per word of Z.

For machines with suitable single precision and no half-word integers (e.g., CDC, Burroughs), REAL*8 and INTEGER*2 are replaced by REAL and INTEGER respectively, and the word-size indicators all take the value 1. The value of EPS should then correspond to *single precision*.

IX.2. Subroutine HASH

This routine is used for storing and searching a table of row names during input of the MPS file. Some machine-dependent "hash function" is required for converting any given 8-character name to a non-negative integer.

Specification

```
        SUBROUTINE HASH( LEN,NEN,NCOLL,
      1 KEY1,KEY2,MODE,KEYTAB,NAME1,NAME2,KA,FOUND )
        INTEGER    KEYTAB(LEN),NAME1(NEN),NAME2(NEN)
        LOGICAL    FOUND
```

Method

See Brent (1973).

Remarks

1. The first application of the required hash function is to the integer parameters KEY1 and KEY2. These variables contain the left and right halves of an 8-character row name, read under 2A4 format. Thus each integer contains 4 characters left-justified. If the resulting bit pattern for each is a valid integer (as is the case on the IBM 360 and 370) then the statement

$$KEY = IABS(IABS(KEY1) - IABS(KEY2))$$

defines a transformation which will produce a suitable non-negative integer KEY from the quantities KEY1 and KEY2. (Note that subtraction of non-negative integers cannot cause overflow.)

2. On some machines (e.g., the Burroughs B6700), reading the integer variable KEY1 (say) under A4 format will set bits which are reserved for the exponent in REAL variables. In such cases the word-length will be greater than 4 characters, and the simplest solution is to perform a right-shift of N bits, where N is chosen to clear the exponent field, e.g.,

$$K1 = IRSHFT(KEY1,N)$$
$$K2 = IRSHFT(KEY2,N)$$
$$KEY = IABS(K1-K2)$$

3. Note that KEY1 and KEY2 are later stored <u>without</u> transforma-
   tion in the integer arrays NAME1 and NAME2. It must be
   possible to perform a test for equality of the form

   IF (KEY1.EQ.NAME1(KT) .AND. KEY2.EQ.NAME2(KT)) GO TO 60

   even if the bit pattern of any of the variables involved is
   not strictly that of an integer. (On the Burroughs B6700 it is
   safer to replace .EQ. by .IS.)

4. The same hash function must be used later within subroutine
   HASH to produce a non-negative integer KEY from the integer
   variables NAME1(KT) and NAME2(KT).

IX.3. <u>Files 5, 6 and 8</u>

   The following file numbers are to some extent built in to
the system:

1. File 5 -- the card reader.

   (a) The SPECS file is normally assumed to be file 5. It
       may be altered by changing the variable ISPECS in
       subroutine GO.

   (b) The MPS file default value is 5. It may be altered at
       run time by a SPECS card of the form

                    INPUT FILE   k

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2. File 6 -- the printer.

All printed output is written on file 6. This includes list-
ings of the SPECS and MPS data, parameters, iteration log,
solution, error messages, etc. If necessary, the characters
"WRITE(6" may be altered in all relevant subroutines.

3. File 8 -- the scratch file.

   (a) This file is used during input of the SPECS file. It
       facilitates interpreting of free-format data in a way
       that is machine-independent.

   (b) The scratch file is later used to store ROW and COLUMN
       names from the MPS data.

   (c) For both purposes the recording format should be fixed
       and blocked, with a logical record length of 8 characters
       (n.b. not 80). The block-size should be at least 800
       or 1600 characters for efficiency.

   (d) The file number may be altered by changing the variable
       ISCRCH in subroutine GO. It must be different from the
       value in variable ISPECS.

   Note that certain files are subject to a Fortran REWIND,
   except if their file numbers are 5 or 6 as the case may be.

IX.4. Using MINOS as a subroutine

As currently implemented, MINOS is intended for use primarily
as a "stand-alone" system, which simply solves a sequence of problems
and then terminates. Figure 5 illustrates the subroutine hierarchy.

Figure 5. Subroutine hierarchy

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In particular,

1. The MAIN program provides a single array of storage to the system, thus:

```
REAL*8      Z(20000)
DATA    NWCORE/20000/

CALL GO( Z,NWCORE )
RETURN
END
```

For solution of larger problems, storage can be easily increased by changing the first two cards. Alternatively an assembly language MAIN program may be used to acquire storage at run-time. (For example, the IBM 370 version of MINOS uses an assembly routine to access all available core via a GETMAIN operation, returns 40K bytes for buffer storage, and then calls subroutine GO as above.)

2. **Subroutine** GO is a control routine at the root of the subroutine tree structure. It may be expanded for specific applications. As given, its main purpose is to call subroutine MINOS and to determine when all of a possible sequence of problems has been processed.

3. Subroutine MINOS and its sub-programs communicate with the outside world by means of certain files, including the SPECS, MPS, BASIS and SOLUTION files. Subroutine MINOS also has output parameters which define:

    (a) the stopping condition;

    (b) the dimensions of the problem that has just been processed;

    (c) the positions within array Z (above) where certain sub-arrays are stored. The latter contain the final values for the solution vector x, the dual variables π, and the state vector defining whether each variable x_j is basic, supe   or nonbasic at upper or lower bound.

The subroutine structure is sufficiently modular to allow use of MINOS in several alternative environments. The principal routine to be modified for a particular application is subroutine GO. It is envisaged that in all cases, run-time parameters and the constraint information will be input via the SPECS and MPS files (rather than by construction in-core).

As an example, the dotted lines in Figure 5 indicate how Matrix Generator and Report Writer modules could be incorporated. The sequence of actions

(a) Generate SPECS file and MPS file

(b) Solve

(c) Write report

could be repeated 1 or more times.

Alternatively it may be convenient to run the Matrix Generator as a separate program. This could generate a series of problem specifications on one SPECS file (since the SPECS file is not rewound after processing of the END card). The corresponding MPS data could be generated as one MPS file (with multiple RHS, RANGE or BOUND sets), or as several separate files.

The Report Writer may access solutions either externally (through the SOLUTION file) or in-core (through the parameter list of subroutine MINOS as noted above). In some cases the SCRATCH file may be useful to provide the ROW and COLUMN names as a sequence of 8-byte records.

A slightly different situation arises if MINOS is to be used iteratively for solving a sequence of linearly constrained problems in which the solution to one problem defines the constraints for the next. This is a possible approach to solving problems with nonlinear constraints. In such cases it is likely that a large portion of the constraint matrix will in fact be constant. The Matrix Generator could set this up as a permanent file, while generating the MPS and SPECS files required for the first Solve step. After each Solve, the "Report Writer" could then

(a) Access the solution as described above, test for convergence and stop if necessary.

(b) Generate new matrix coefficients and merge with the permanent file to produce a new MPS file.

(c) Generate an appropriately modified SPECS file.

(d) Modify relevant parameters used in the objective function.

(e) Return to the Solve step.

Although purely in-core operation is generally preferable, most
of the computation for large problems will occur in the Solve step.
The interspersed file-handling should therefore incur relatively
negligible cost.

X. System Information -- IBM Version

This section illustrates the Job Control required to compile, run and modify MINOS on an IBM System 360 or 370 operating under OS. Some of the parameters used may be peculiar to OS/VS2 and/or the particular installation at the Stanford Linear Accelerator Center. Various DSNAMES and VOLUME names may also need to be altered. However, the main requirements should be apparent.

X.1. To Compile MINOS and Save the Resulting Load Module

In this example an executable program called NLPPGM is created as a member of the Partitioned Data Set called WYL.W8.MXS.LOAD, which will reside on a 2314 disk unit called WYLOOA.

The compile time for 8000 Fortran source cards using the IBM Fortran (H Extended) compiler with full optimization is approximately 40 seconds on an IBM 370/168.

```
//   JOB   ,TIME=(1,00),REGION=256K
//   EXEC   FORTHCL,PARM.FORT='OPT=2,NOMAP,NOSOURCE',
//          FORTPRT=DUMMY,
//          FORTTIM='(1,00)',FORTUER=NEW
//FORT.SYSIN  DD  *


      Fortran source


/*
//LKED.SYSLMOD  DD  DSN=WYL.W8.MXS.LOAD,UNIT=DISK,
//          DISP=(NEW,CATLG),VOL=SER=WYLOOA,
//          DCB=(RECFM=U,LRECL=3520,BLKSIZE=3520),
//          SPACE=(TRK,(40,5,1),,CLSE)
//LKED.SYSIN  DD  *
  ENTRY   MAIN
  NAME    NLPPGM
/*
```

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X.2. To Run MINOS

The following job is suitable for running the load module created by the preceding example, or by the example in Section X.4. below. Note that the DDNAMES of the form 'FTnnF001' do not have the usual 'GO.' attached, because the EXEC card specifies NLPPGM directly rather than through a Cataloged Procedure.

- The 3rd and 4th characters in each 'FTnnF001' DDNAME refer to file numbers which are used in Fortran I/O statements such as READ(nn, f).

FT05F001   Refers to the card reader.

FT06F001   Refers to the printer.

FT08F001   Defines the scratch file. Note that the logical record length is LRECL=8, not LRECL=80.

FT09F001   etc. refer to various basis and data files as required for a particular run. Note that DISP=SHR implies that a file already exists (as does DISP=OLD). If necessary some dummy file should be set up before the run. Otherwise, DISP=SHR must be replaced by DISP=(NEW,KEEP) and appropriate DCB and SPACE parameters must be defined.

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```
//   JOB  ,TIME=(0,15),REGION=320K,MSGLEVEL=(0,0)
//MINOS  EXEC  PGM=NLPPGM
//STEPLIB    DD   DSN=WVL.W8.MXS.LOAD,DISP=SHR
//FT06F001   DD   SYSOUT=H,
//           DCB=(RECFM=FBA,LRECL=133,BLKSIZE=3458,BUFNO=2)
//FT08F001   DD   UNIT=SYSDA,SPACE=(TRK,(4,4)),
//           DCB=(RECFM=FB,LRECL=8,BLKSIZE=1600)
//FT09F001   DD   UNIT=DISK,VOL=SER=SCR001,DSN=WVL.W8.MXS.BASIS1,
//           DISP=SHR
//FT10F001   DD   UNIT=DISK,VOL=SER=SCR001,DSN=WVL.W8.MXS.BASIS2,
//           DISP=SHR
//FT11F001   DD   UNIT=DISK,VOL=SER=SCR001,DSN=WVL.W8.MXS.MPSDATA
//           DISP=SHR
//FT05F001   DD   *


    SPECS

    MPS data, if INPUT FILE 11 is not used  (FT11F001  above)

    Any data required by subroutine CALCFG


/*
```

X.3. To Compile CALCFG and Other Subroutines and Link to MINOS

The following job shows how to link one or more new or altered
subroutines to MINOS and run the resulting program. This mode of
operation is normally used to include a user's own nonlinear objec-
tive subroutine (CALCFG) or to test special versions of some of
the subroutines in MINOS.

In the Link Edit step, the two XLIB cards refer to an
existing load module. The new load module is not saved.

```
//    JOB   ,TIME=(0,15),REGION=320K,MSGLEVEL=(0,0)
//CLG  EXEC   FORTHCLG,PARM.FORT='OPT=2,NOMAP,NOSOURCE',
//            FORTPRT=DUMMY,LKEDPRT=DUMMY,
//            FORTTIM='(0,40)',FORTUER=NEW
//FORT.SYSIN  DD  *


       Fortran source for new subroutines


/*
//LKED.XLIB    DD   DSN=WVL.W8.MXS.LOAD,DISP=SHR
//LKED.SYSIN   DD  *
  INCLUDE XLIB(NLPPGM)
  ENTRY   MAIN
/*
//GO.FT08F001    DD   UNIT=SYSDA,SPACE=(TRK,(4,4)),
//               DCB=(RECFM=FB,LRECL=8,BLKSIZE=1600)
//GO.FT09F001    DD   UNIT=DISK,VOL=SER=SCR001,DSN=WVL.W8.MXS.BASIS1,
//               DISP=SHR
//GO.FT10F001    DD   UNIT=DISK,VOL=SER=SCR001,DSN=WVL.W8.MXS.BASIS2,
//               DISP=SHR
//GO.FT11F001    DD   UNIT=DISK,VOL=SER=SCR001,DSN=WVL.W8.MXS.MPSDATA,
//               DISP=SHR
//GO.SYSIN  DD  *


    SPECS

    MPS data, if INPUT FILE 11 is not used  (GO.FT11F001  above)

    Any data required by subroutine CALCFG


/*
```

X.4. To Create and Save a New Load Module

       In this example the IKED.SYSIMOD card specifies where the

new version of MINOS is to be saved. About 40 tracks of 2314 disk

are required, or about 20 tracks of 3330 disk.

If necessary, any object decks (e.g., from assembly language versions of certain subroutines) should be included immediately after the IKED.SYSIN card.

```
//  JOB  ,TIME=(0,15),REGION=256K,MSGLEVEL=(0,0)
//CL   EXEC  FORTHCL,PARM.FORT='OPT=2,NOMAP,NOSOURCE',
//           FORTPRT=DUMMY,
//           LKEDPRT=DUMMY,
//           FORTTIM='(0,40)',FORTVER=NEW
//FORT.SYSIN  DD  *


       Fortran source for new subroutines


/*
//LKED.SYSLMOD  DD  DSN=WYL.W8.MXS.LOAD,UNIT=DISK,
//           DISP=(NEW,KEEP),VOL=SER=WYLOOB,
//           DCB=(RECFM=U,LRECL=3520,BLKSIZE=3520),
//           SPACE=(TRK,(40,5,1),RLSE)
//LKED.XLIB   DD  DSN=WYL.W8.MXS.LOAD,DISP=SHR
//LKED.SYSIN  DD  *
  INCLUDE XLIB(NLPPGM)
  ENTRY   MAIN
  NAME    NLPPGM
/*
```

Note that:

--the old load module library (WYL.W8.MXS.LOAD) is assumed to be CATALOGed on the appropriate volume (e.g., WYLOOA as in the first job);

--the new load module library has the same name and must therefore reside on a different volume, in this case WYLOOB;

--the name of the new executable program is also the same as

before (NLPPGM).

The purpose here is to make creation of the new version of

MINOS transparent to the previous two jobs, which use whatever

version happens to be referred to by the system CATALOG. It should

now be possible to type a command of the form

RECATALOG   LOAD   ON   WYLOOB

on the local terminal system, and the deed will be done.

If the required terminal system facilities are not available,

the following job illustrates how the RECATALOG operation may be

accomplished using JCL alone. The point here is that the new

load module can be cataloged only if the old version is uncataloged

in a previous job step.

```
//   JOB   ,TIME=(0,15),REGION=256K,MSGLEVEL=(0,0)
//CL    EXEC   FORTHCL,PARM.FORT='OPT=2,NOMAP,NOSOURCE',
//             FORTPRT=DUMMY,
//             LKEDPRT=DUMMY,
//             FORTTIM='(0,40)',FORTVER=NEW
//FORT.UNCAT  DD  DSN=WYL.W8.MXS.LOAD,DISP=(SHR,UNCATLG)
//FORT.SYSIN  DD  *
```

Fortran source for new subroutine:

```
/*
//LKED.SYSLMOD  DD  DSN=WYL.W8.MXS.LOAD,UNIT=DISK,
//         DISP=(NEW,CATLG),VOL=SER=WYL00B,
//         DCB=(RECFM=U,LRECL=3520,BLKSIZE=3520),
//         SPACE=(TRK,(40,5,1),RLSE)
//LKED.XLIB   DD  DSN=WYL.W8.MXS.LOAD,DISP=SHR,
//             UNIT=DISK,VOL=REF=*.CL.FORT.UNCAT
//LKED.SYSIN  DD  *
  INCLUDE XLIB(NLPPGM)
  ENTRY   MAIN
  NAME    NLPPGM
/*
```

Beware that the old load module will be uncataloged even if the
compile or link edit steps fail.

XI. Further Examples

XI.1. Example 2 -- An Unconstrained Optimization Problem

The following data illustrates minimization of the well known function (Rosenbrock, 1960)

$$f(x) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

without any general linear constraints.

Notes

1. At least one "constraint" is required in all problems. In this case the simplest technique is to include a free (non-binding) row, called DUMMYROW here. The basis matrix will remain $B = 1$ throughout, corresponding to the slack variable on the free row.

2. The COLUMNS section contains just a list of the variable names, and the RHS section may be null.

3. The bounds $-10 \leq x_j \leq 10$ are included via the bound set BNDX, mainly as a matter of good practice. (Their presence does not imply additional work.) A uniform set of this kind could also be requested by the SPECS cards

        LOWER BOUND   -10.0

        UPPER BOUND   10.0

FUNCTION AND GRADIENT CALCULATION (for example 2)

```
      SUBROUTINE CALCFG( MODE,N,X,F,G,NSTATE,NPROB )
      IMPLICIT    REAL*8(A-H,O-Z)
      REAL*8      X(N),F,G(N)
C
C     ------------------------------------
C     ROSENBROCK'S BANANA FUNCTION
C     ------------------------------------
      S = X(2) - X(1)**2
      T = 1.0  - X(1)
      F = 100.0*S**2 + T**2
      G(1) = -400.0*S*X(1) - 2.0*T
      G(2) =  200.0*S
      RETURN
C
C     END OF CALCFG
      END
```

SPECS

```
BEGIN ROSENBROCK

   ITERATIONS              100

   NONLINEAR VARIABLES       2
   SUPERBASICS               2
   HESSIAN DIMENSION         2

   LINESEARCH TOLERANCE    0.1
END
```

MPS DATA

```
NAME           ROSENBROCK
ROWS
 N  DUMMYROW
COLUMNS
    X1
    X2
RHS
BOUNDS
 LO BNDX       X1            -10.0
 UP BNDX       X1             10.0
 LO BNDX       X2            -10.0
 UP BNDX       X2             10.0
 FX INITIAL    X1             -1.2
 FX INITIAL    X2              1.0
ENDATA
```

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4. The INITIAL bound set illustrates how the starting point $(x_1, x_2) = (-1.2, 1.0)$ may be specified. Since the limit on SUPERBASICS is sufficiently high (2 here) both x_1 and x_2 will initially be superbasic at the specified values.

5. If the INITIAL bound set were not used, x_1 and x_2 would initially be nonbasic at the smaller of their bounds (in absolute magnitude), in this case the upper bounds 10.0.

6. If the user wished the variables to be genuinely unconstrained, the following bounds would have to be used:

```
BOUNDS
    FR FREE       X1
    FR FREE       X2
    FX INITIAL    X1              -1.2
    FX INITIAL    X2               1.0
```

where the INITIAL bounds are again optional. If the INITIAL bounds were absent, x_1 and x_2 would both be regarded initially as nonbasic at value zero.

XI.2. Example 3 -- A Quadratic Program

The following data illustrates solution of the quadratic program

$$\text{minimize} \quad \frac{1}{2} x^T Q x + c^T x$$

$$\text{subject to } Ax \leq b, \qquad x \geq 0$$

FUNCTION AND GRADIENT CALCULATION   (for example 3)

```
      SUBROUTINE CALCFG( MODE,N,X,F,G,NSTATE,NPROB )
      IMPLICIT    REAL*8(A-H,O-Z)
      REAL*8      X(N),F,G(N)
      COMMON      Q(3,3)
C
C     COMPUTATION OF  F = 1/2 X'*Q*X,     G = Q*X     FOR GIVEN Q.
C     THE   COMMON   STATEMENT AND   SUBROUTINE SETQ   ARE PROBLEM
C     DEPENDENT.
C
C
      IF (NSTATE.EQ.1) CALL SETQ( Q,N )
C
      F = 0.0
      DO 300 I = 1, N
         S = 0.0
         DO 200 J = 1, N
            S = Q(I,J)*X(J) + S
  200    CONTINUE
         F    = X(I)*S + F
         G(I) = S
  300 CONTINUE
C
      F = 0.5*F
      RETURN
      END
```


SPECS

```
BEGIN QP
   ITERATIONS              50
   NONLINEAR VARIABLES      3
END
```


MPS DATA

```
NAME           QP
ROWS
 N C
 L A
COLUMNS
    X1         C         -8.0         A              1.0
    X2         C         -6.0         A              1.0
    X3         C         -4.0         A              2.0
RHS
    B          A          3.0
ENDATA
```

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for the particular data

$$Q = \begin{bmatrix} 4 & 2 & 2 \\ 2 & 4 & 0 \\ 2 & 0 & 2 \end{bmatrix} \qquad c = \begin{bmatrix} -8 \\ -6 \\ -4 \end{bmatrix}$$

$$A = \begin{bmatrix} 1 & 1 & 2 \end{bmatrix} \qquad b = 3$$

Notes

1. The vector c is included as a free row in the constraint matrix, as in standard LP.

2. The $\leq$ shown above may of course include all forms of inequality as desired.

3. Often Q will not involve all variables. In such cases it is preferable to order the variables so that the nonzero rows and columns of Q are grouped together thus:

$$Q = \begin{bmatrix} \bar{Q} & \\ & 0 \end{bmatrix}$$

The number of NONLINEAR VARIABLES and the parameter N of subroutine CALCFG will then be the dimension of $\bar{Q}$.

4. It is assumed that Q is initialized by some subroutine SETQ during the first entry to subroutine CALCFG. The COMMON statement is also problem-dependent (and is included to ensure that Q retains its values for later entries). Otherwise subroutine CALCFG is suitable for any quadratic function involving the first N components of x.

5. Beware -- if $c \neq 0$, the factor $1/2$ makes a vital difference to the function being optimized.

6. The optimal solution is

$$x = (1.3333, \ 0.7777, \ 0.4444)$$

$$\tfrac{1}{2}x^T Q x = 8.2222$$

$$c^T x = -17.1111$$

$$\text{objective} = -8.8888$$

XI.3.   Example 4 -- Linear Least Squares

Data-fitting can give rise to a constrained least squares problem of the form

$$\underset{x}{\text{minimize}} \quad \|Xx - y\|_2$$

$$\text{subject to } Ax \geq b, \qquad \ell \leq x \leq u.$$

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This may be solved with MINOS by expressing the problem as

$$\underset{r,x}{\text{minimize}} \quad \frac{1}{2} r^T r$$

$$\text{subject to} \quad \begin{bmatrix} & A \\ I & X \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} \begin{matrix} \geq \\ = \end{matrix} \begin{bmatrix} b \\ y \end{bmatrix}$$

$$r \quad \text{"free,"} \qquad \ell \leq x \leq u \, .$$

Notes

1. As usual, $Ax \geq b$ may include all types of inequality.

2. $r = y - Xx$ is the "residual vector," and $r^T r$ is the "sum of squares."

3. The objective function and gradient are easily programmed as $f(r) = 1/2 \ r^T r, \ g(r) = r.$

4. Use of MINOS in this context is certainly not claimed to be the most stable or efficient method known. Stability depends on the condition of a certain submatrix extracted from

$$\begin{bmatrix} A \\ X \end{bmatrix} \, .$$

If it is known that most of the variables x will be away from their bounds, efficiency could be improved by using the option

(say). This would allow up to 10 variables at a time to be
added to the set currently being optimized, instead of the
usual 1.

XI.4.   <u>Example 5 -- The Discrete</u> ℓ_1 <u>Problem</u>

An apparently similar data-fitting problem is

$$\underset{x}{\text{minimize}} \quad \|Xx - y\|_1$$

$$\text{subject to} \quad Ax \geq b$$

where $\|r\|_1 \equiv \sum_i |r_i|$. However, this is best solved by means of
the following purely <u>linear</u> program:

$$\underset{\lambda,\mu}{\text{maximize}} \quad [y^T \quad b^T] \begin{bmatrix} \lambda \\ \mu \end{bmatrix}$$

$$\text{subject to} \quad [X^T \quad A^T] \begin{bmatrix} \lambda \\ \mu \end{bmatrix} = 0 ,$$

$$-1 \leq \lambda \leq 1, \quad \mu \geq 0 .$$

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<u>Notes</u>

1. The solution x is recovered as the "simplex multipliers."

2. The optimal value of $\|Xx - y\|_1$ is the sum of the absolute values of the "reduced costs" associated with λ. (It is also the maximal value of $y^T\lambda + b^T\mu$.)

3. If a particular row in A, b is required to be an equality constraint, the corresponding component of μ should be a free variable.

4. It does not appear simple to include the bounds $\ell \leq x \leq u$ except as part of $Ax \geq b$. In such cases it may be best to solve the original problem directly as a linear program, thus:

$$\begin{array}{c} \text{minimize} \\ r,s \end{array} \quad \begin{bmatrix} e^T & e^T \end{bmatrix} \begin{bmatrix} r \\ s \end{bmatrix}$$

$$\text{subject to} \quad \begin{bmatrix} & & A \\ & & \\ I & -I & X \end{bmatrix} \begin{bmatrix} r \\ s \\ x \end{bmatrix} \begin{array}{c} \geq \\ \\ = \end{array} \begin{bmatrix} b \\ \\ y \end{bmatrix},$$

$$r \geq 0, \qquad s \geq 0, \qquad \ell \leq x \leq u,$$

where $e^T = [\; 1 \; 1 \; \cdots \; 1 \;]$.

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XII. Conclusion

XII.1. A Word to the Linear Programmer

Those familiar with commercial mathematical programming systems will have noticed an absence of the usual algorithmic options (dual simplex method, GUB, parametrics, etc.), not to mention the normal wealth of data revision and file handling capabilities. However, it should be clear that MINOS is intended for use as an extension of such systems, not as a substitute. Certainly the necessary constraint data can be prepared and revised using standard matrix generators and the above-mentioned facilities. In fact, we would recommend that any sizeable model be thoroughly "debugged" as a purely linear program, using either MINOS or (preferably) a commercial system. An appropriate set of bounds can be used to fix the nonlinear variables at plausible values during this debugging stage. The final optimal solution thus obtained will generally provide a good starting point when the nonlinear objective is introduced and the bounds on nonlinear variables are relaxed. (In this context it is essential to follow the suggestions of Section VI.)

Note that if the constraint data has already been debugged, the INITIAL bounds set of Section III.2 allows this fixing and subsequent freeing of nonlinear variables to be accomplished automatically within a single run.

XII.2. To the Nonlinear Programmer

Nonlinear programming might reasonably be expected to include nonlinear constraints--another vital subject which has received no mention here. In answer to this point we can only say that there is something fundamentally special about linear constraints, namely that it is possible to stay _on_ them with relative ease. In implementing MINOS we have shown that a significant class of large-scale linearly constrained problems can be solved with only moderately more effort than would be required by a linear program of the same dimensions. This situation changes dramatically with the slightest mention of nonlinear constraints (unless perhaps they are included as penalty terms in the objective--a strategy which normally leads to ill-conditioned problems and/or crude approximations to the solution).

To illustrate, the Generalized Reduced-Gradient method of Abadie and Carpentier (1969) is directly analogous to the reduced-gradient approach used in MINOS, but an order of magnitude more work is involved (for the more general problem) in maintaining feasibility with respect to the nonlinear constraints. Recently, a derivative of MINOS has been developed (MINOS/GRG, Jain, Lasdon and Saunders, 1976) with the intention of handling about 100 nonlinear constraints and several hundred sparse linear constraints. As might be expected, run times on significantly nonlinear problems have proved to be 5 or 10 times what would be typical if all constraints were linear. This is not to discourage use of the system on real problems, but it should

serve to underscore the quantum difference between the two types of nonlinear program. Alternative implementations must be pursued to help narrow the gap.

In the meantime, we hope that for many users, the facilities now available in MINOS will provide a useful first step into the nonlinear world.

Acknowledgements

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SECURITY CLASSIFICATION OF THIS PAGE (When Data Entered)

REPORT DOCUMENTATION PAGE		READ INSTRUCTIONS BEFORE COMPLETING FORM
1. REPORT NUMBER SOL 77-9	2. GOVT ACCESSION NO.	3. RECIPIENT'S CATALOG NUMBER
4. TITLE (and Subtitle) MINOS: A Large-Scale Nonlinear Programming System (For Problems with Linear Constraints) User's Guide		5. TYPE OF REPORT & PERIOD COVERED Technical Report
		6. PERFORMING ORG. REPORT NUMBER
7. AUTHOR(s) Bruce A. Murtagh and Michael A. Saunders		8. CONTRACT OR GRANT NUMBER(s) N00014-75-C-0865
9. PERFORMING ORGANIZATION NAME AND ADDRESS Department of Operations Research - SOL Stanford University Stanford, CA 94305		10. PROGRAM ELEMENT, PROJECT, TASK AREA & WORK UNIT NUMBERS NR-047-143
11. CONTROLLING OFFICE NAME AND ADDRESS Operations Research Program, Code 434 Office of Naval Research Arlington, VA 22217		12. REPORT DATE February
		13. NUMBER OF PAGES 127
14. MONITORING AGENCY NAME & ADDRESS(if different from Controlling Office)		15. SECURITY CLASS. (of this report) UNCLASSIFIED
		15a. DECLASSIFICATION/DOWNGRADING SCHEDULE

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18. SUPPLEMENTARY NOTES

19. KEY WORDS (Continue on reverse side if necessary and identify by block number)

Linearly constrained optimization; optimization; linear programming; nonlinear programming; sparse matrices; MINOS; mathematical software; large-scale optimization.

20. ABSTRACT (Continue on reverse side if necessary and identify by block number)
MINOS is a Fortran program designed to minimize a linear or nonlinear function subject to linear constraints, where the constraint matrix is in general assumed to be large and sparse. The User's Guide contains an overview of the MINOS System, including descriptions of the theoretical algorithms as well as the details of implementation. The Guide also provides complete instructions for the use of MINOS, and illustrates the diversity of application by several examples.

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